OS Scheduling/Buffering

- · Scheduling
 - Kernel/user space
 - Processes/threads
 - Interrupts
 - Non-preemptive/Preemptive
 - Realtime
- · Buffering
 - Block based
 - Shared memory/Memory mapping
 - Network smoothing
 - Playback buffers

Kernel/user space Changing from Kernel to user mode called a Restricted User process "context change" System calls Supervisor Kernel Drivers

Processes/Threads

- · Process
 - a task being run (often simultaneously) with other processes
 - Must take turns on the CPU (timeslicing) unless there are multiple CPUs (multiprocessing processes may be allocated to CPUs)
 - Managed by the OS (scheduled)
 - Kept apart to reduce chance of system failure (but may communicate by IPC Inter Process Communication)
- · Threads
 - Process splits itself into two or more simultaneously running tasks
 - Share process state (are more dependent)
 - Share memory directly
 - Faster context switching

Interrupts

- · Signal (event) from a device to the processor
- Causes the processor to stop what it is doing and execute an "interrupt handler" (context switch)
- Can be from audio device, timer...

access

access

- Timer interrupt increments a clock, used by the kernel to switch processes
- Can be disabled (except "non-maskable interrupts") to stop interrupts interrupting other
- Interrupt routine should be fast (another interrupt could be on the way)

(Non-)preemptive scheduling

- · Non-preemptive multitasking
 - OS does not try to guess when a process has finished (does not pre-empt it)
- · Pre-emptive multitasking
 - OS "interrupts" a process (even if it is not complete) and gives control to another process
 - Hardware interrupts can also pre-empt a process
 - Eg Windows, Linux...

Realtime Scheduling

- A Real Time Operating System (RTOS), is an operating system that has been developed for real-time applications.
- Typically used for embedded applications they usually have the following characteristics:
 - Small footprint (doesn't use much memory)
 - Pre-emptable (any hardware event can cause a task to run)
 - Multi-architecture (code ports to another type of CPU)
 - Many have predictable response-times to electronic events
- Minimise interrupt disable period
- Application can request delay, jitter bounds. OS schedules processes according to requests, to meet application requirements.

Buffering

- · Block based
 - Device/driver handles data in blocks
 - more efficient than, eg, single bytes: less interrupts/OS calls
 - Bigger blocks = more efficient but more delay
 - Smaller blocks = less efficient but less delay
 - Codecs sometimes restrict choice (eg lpc, mp3...)

Shared memory/Memory mapping

- · Shared memory
 - Processes/threads communicate/exchange data via memory accessible to all processes
 - Processes have to manage access (semaphores)
- Memory mapping
 - An area of memory used by the kernel is mapped into user space (user process can access directly)
 - Eg Audio buffer, Framebuffer (video)

Network smoothing

- Variable rate output from a codec (eg Iframes, P-frames, B-frames)
- Constant bitrate output from a network interface
- Network buffers "smooth" the codec output to fit the network capacity
- · Causes some delay

Playout buffers

- Receiver receives data with jitter (from network/OS scheduling...)
- Adding a short delay (holding the data in a buffer) may smooth the jitter
- Buffer size depends on type of application (interactive/non-interactive)
- Interactive buffer size=min("worst case jitter seen so far", 150ms) (may be too conservative)
- Non-interactive buffer size=0.5s (allows for reasonable speed channel hopping)