

Scalable Shape Analysis For Systems Code

Hongseok Yang (Queen Mary, Univ. of London)

Joint work with

Oukseh Lee, Josh Berdine, Cristiano Calcagno, Byron Cook,
Dino Distefano, Peter O'Hearn

Open Question, Summer 2006

Can we automatically prove the pointer safety of programs $\geq 10K$ in separation logic?

Demo

Result

Designed and implemented a shape analyzer for C.

Program	LOC	Sec	MB	Memory leaks	Dereference errors
<code>scull.c</code>	1010	0.21	1.47	1	0
<code>class.c</code>	1983	6.68	8.36	2	1
<code>pci-driver.c</code>	2532	0.79	3.19	0	0
<code>ll_rw_blk.c</code>	5469	997.66	523.22	3	1
<code>cdrom.c</code>	6218	91.88	84.30	0	2
<code>md.c</code>	6635	1440.53	814.45	6	5
<code>t1394Diag.c</code>	10240	145.95	71.27	33	10

Table 1. Experimental Results. Performed on an Intel Core Duo 2GHz with 2GB.

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- Programs of 1010 ~ 10240 LOC.
- Includes the full t1394Diag firewire driver.

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- Found memory leaks and safety errors.
- Fixed those errors. Then, verified the integrity of pointer manipulation.

Results

Automatic Pointer-Safety Prover in Sep. Logic

{emp} t1394Diag_main() {emp}

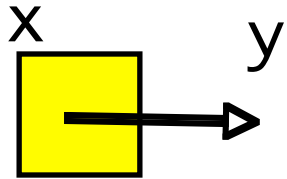
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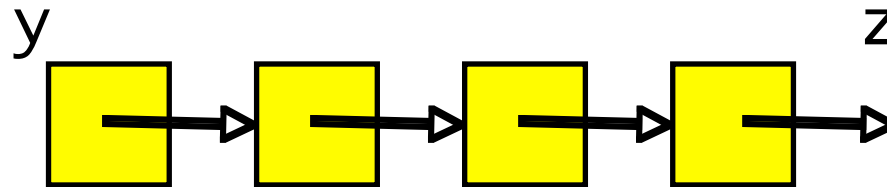
Separation Logic

$x \mapsto y$

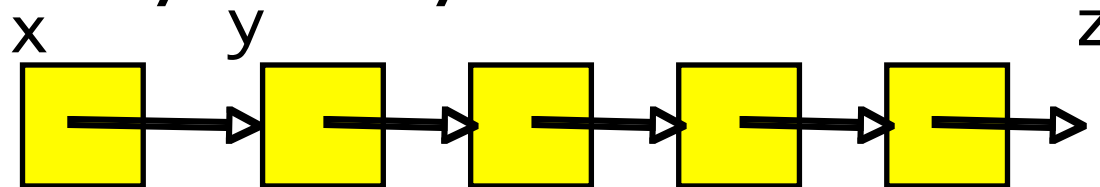


emp

$\text{lsne } y \ z$



$x \mapsto y * \text{lsne } y \ z$



$\exists w', v'. (z=y \wedge w' \neq v') \wedge (x \mapsto w' * \text{lsne } w' \ v' * \text{lsne } y \ v')$

Symbolic Heaps

Separation logic formulas of the form:

$$\exists w', v'. (y=z \wedge w' \neq v') \wedge (x \mapsto w' * \text{lsne } w' \ v' * \text{lsne } y \ v')$$

Symbolic Heaps

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Function “Abs”

$Abs : SH \rightarrow SH$

- Forgets the length of lists. Removes unnecessary primed vars.
- Ensures the termination of the analysis.
- E.g. $z \neq x' \wedge (x \mapsto x' * x' \mapsto 0 * \text{Isne } y \ y' * \text{Isne } y' \ 0)$.

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Shape Analysis

- Run a program symbolically with symbolic heaps.
- Apply abstraction periodically to ensure termination.
- Repeat until the fixpoint is reached.

Interprocedural Shape Analysis

- For each procedure, create a table that records all the past analysis results.
- Table for create():

Precond.	Postcondition
emp	$ret=0 \wedge emp,$ $ret \mapsto 0,$ $!s \text{ ret } 0$
...	...

- Use the table whenever possible.
- The implementation is based on a more sophisticated algorithm (RHS).

```
L create() {...}
```

```
L append(L a,L b) {...}
```

```
void main() {  
  x=create();  
  y=create();  
  x=append(x,y);  
}
```

```
L create() {...}
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```
L append(L a,L b) {...}
```

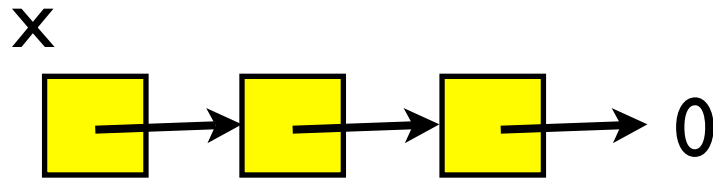
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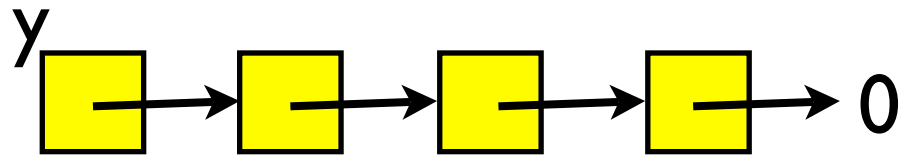
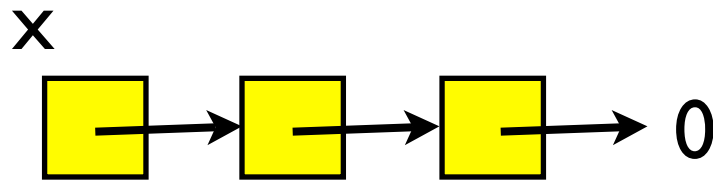
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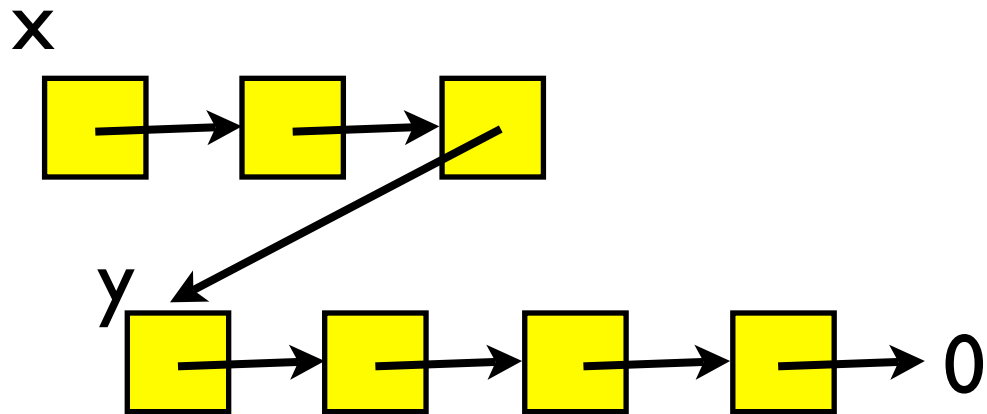
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ret=0 \wedge emp,

ret \mapsto 0,

!sne ret 0

L append(L a,L b) {...}

emp

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void main() {  
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	emp	ret=0 \wedge emp,	ret \mapsto 0,	!sne ret 0
--	-----	---------------------	------------------	------------

L create() {...}

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void main() { emp

x=create(); x=0 \wedge emp

x \mapsto 0

!sne x 0

y=create();
x=append(x,y);
}

emp	ret=0 \wedge emp,	ret \mapsto 0,	!sne ret 0
x=0 \wedge emp			
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x=0 \wedge emp	x=ret=0 \wedge emp,	x=0 \wedge ret \mapsto 0,	x=0 \wedge !sne ret 0
x \mapsto 0	ret=0 \wedge x \mapsto 0,	x \mapsto 0 * ret \mapsto 0,	x \mapsto 0 * !sne ret 0
!sne x 0	ret=0 \wedge !sne x 0,	!sne x 0*ret \mapsto 0,	!sne x 0*!sne ret

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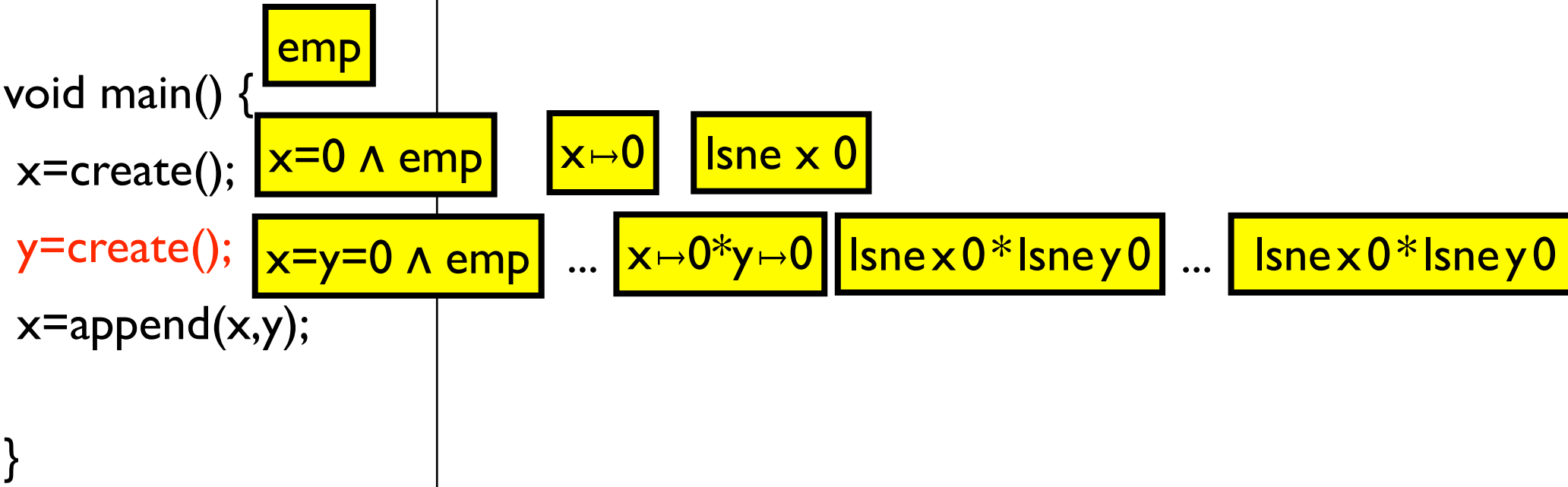
x=0 \wedge emp

x \mapsto 0

!sne x 0

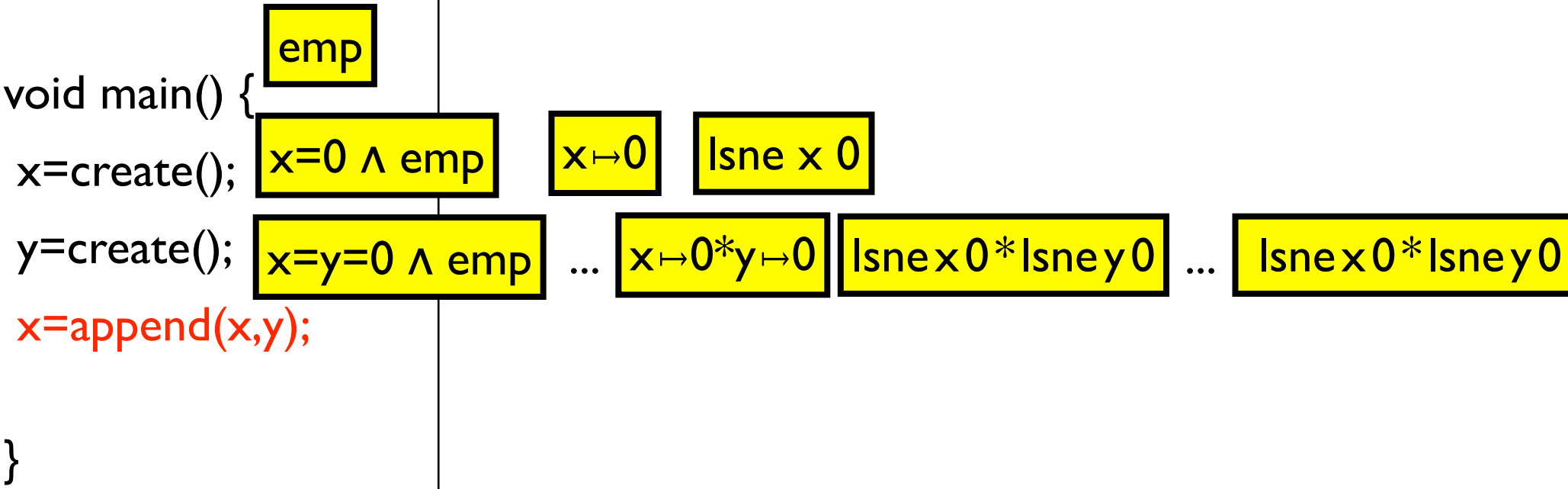
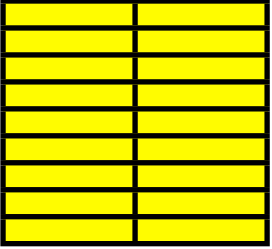
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x=0 \wedge emp	x=ret=0 \wedge emp,	x=0 \wedge ret \mapsto 0,	x=0 \wedge !sne ret 0
x \mapsto 0	ret=0 \wedge x \mapsto 0,	x \mapsto 0 * ret \mapsto 0,	x \mapsto 0 * !sne ret 0
!sne x 0	ret=0 \wedge !sne x 0,	!sne x 0*ret \mapsto 0,	!sne x 0*!sne ret

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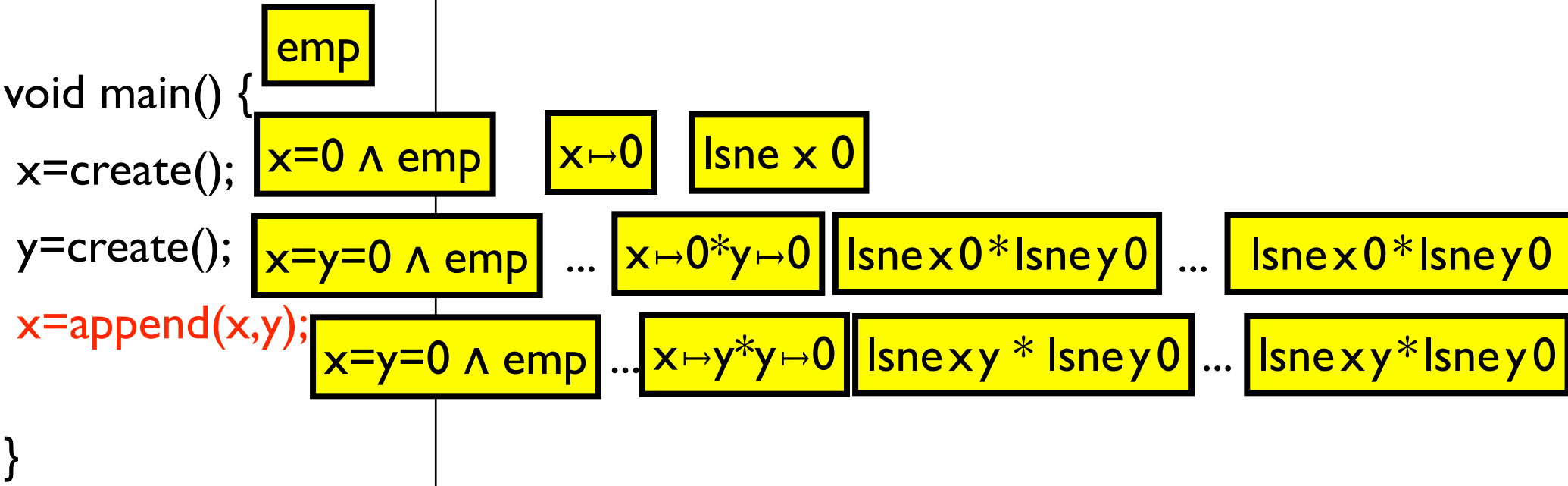
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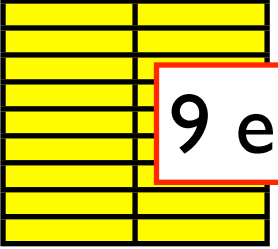
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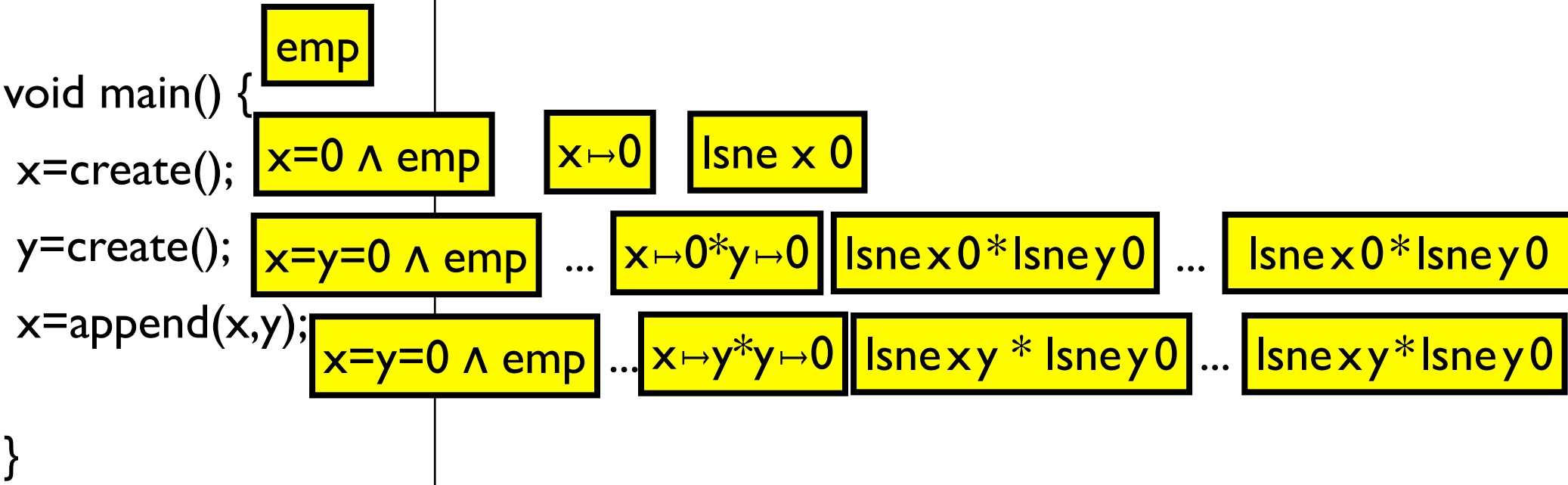
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	x \mapsto 0	ret=0 \wedge x \mapsto 0,	x \mapsto 0 * ret \mapsto 0,	x \mapsto 0 * lsne ret 0
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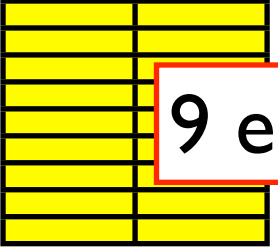
4 entries, 12 results

9 entries, 12 results



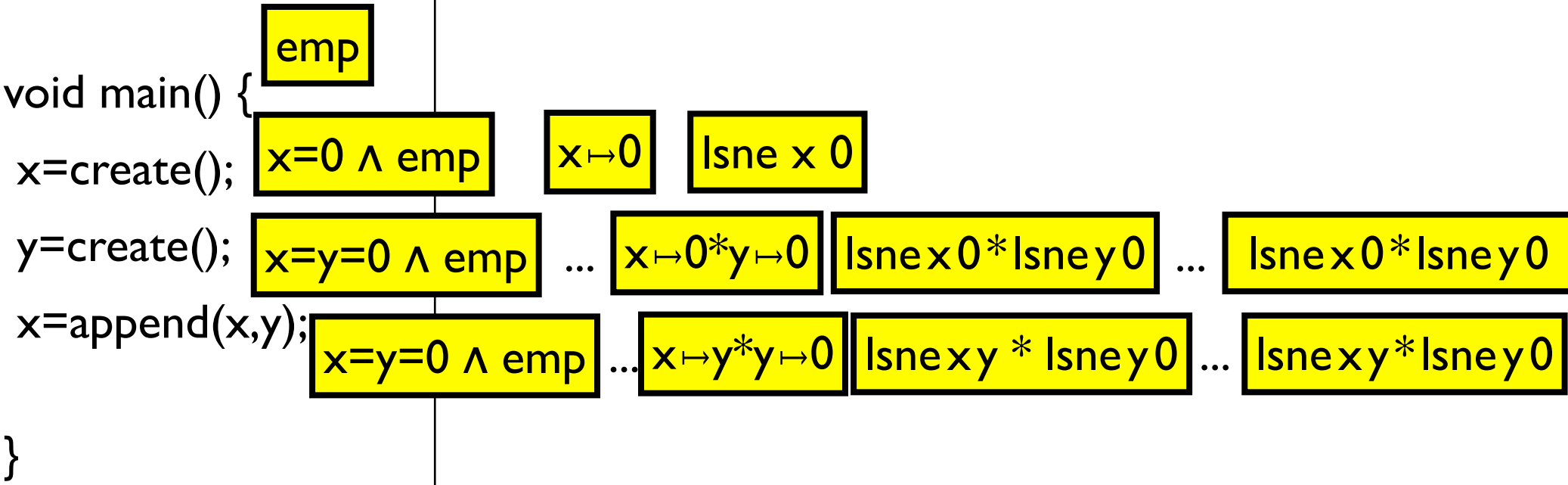
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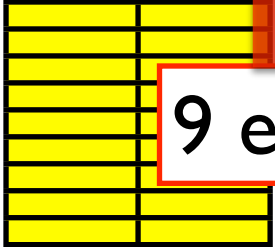


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x \mapsto 0

!sne x 0

y=create();

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... !sne x 0*!sne y 0

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x=y=0 \wedge emp

... x \mapsto y*y \mapsto 0

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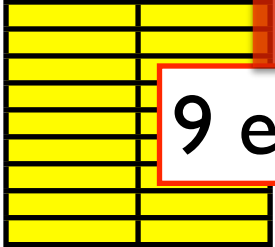
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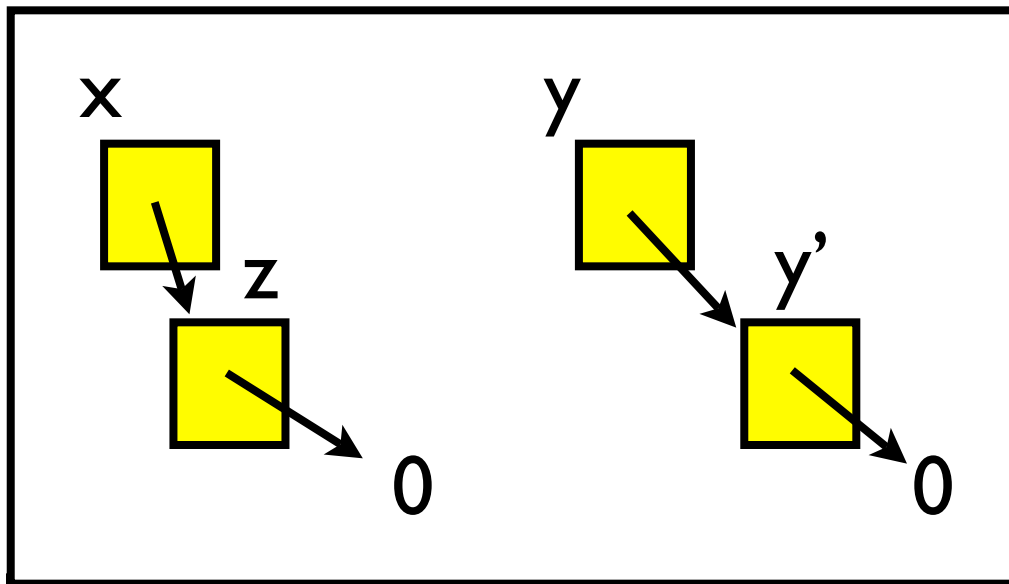
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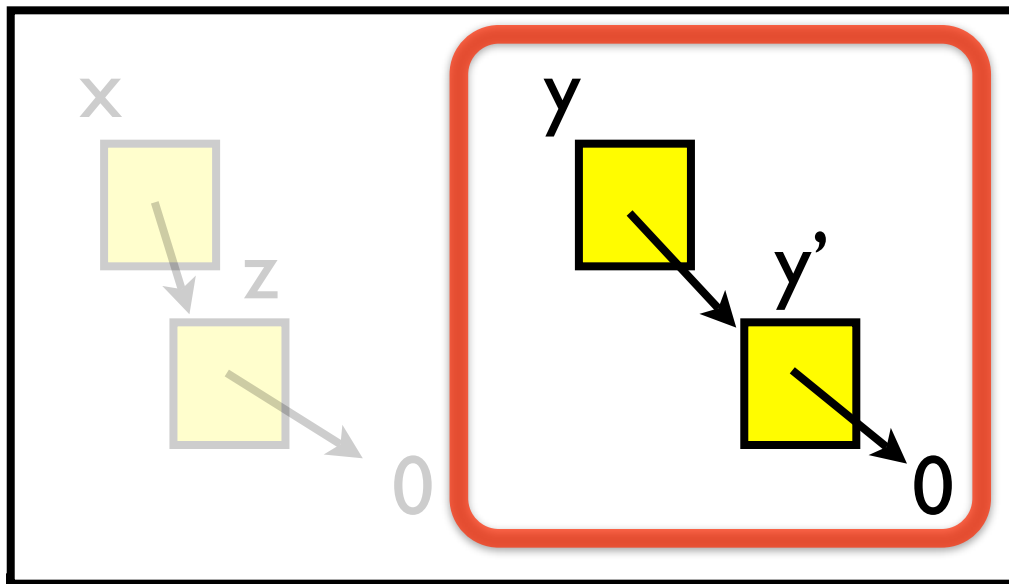
Optimization I: Localization

- Pass & change only the part of a symbolic heap, that is reachable from the parameters. [Rinetzky+POPL05, Gotsman+SAS06.]
- Based on frame rule in separation logic. [CSL01]
- E.g.
 - $\text{lsne } x \ z * z \mapsto 0 * \text{lsne } y \ y' * \text{lsne } y' \ 0, \quad \text{dispose}(y)$



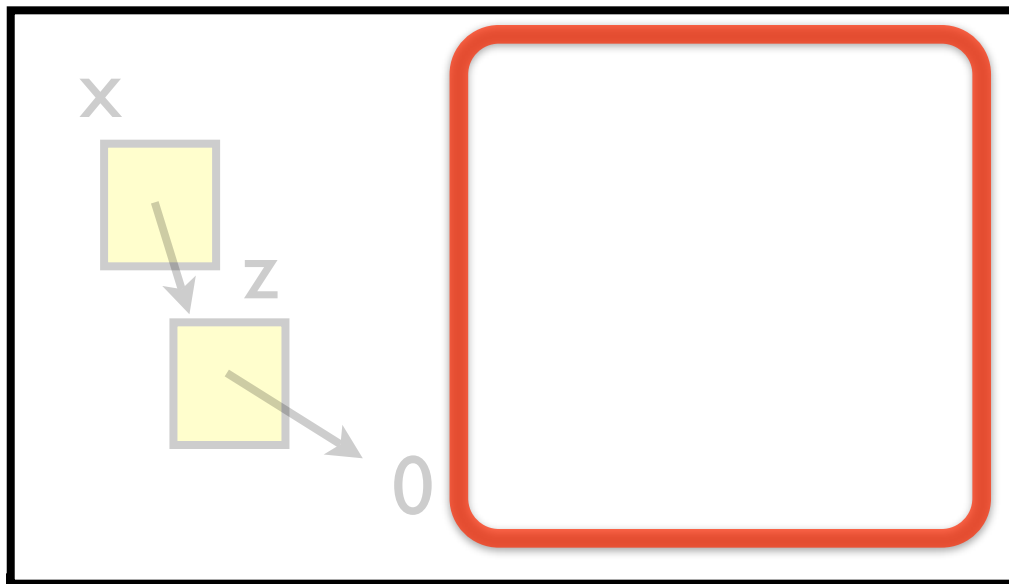
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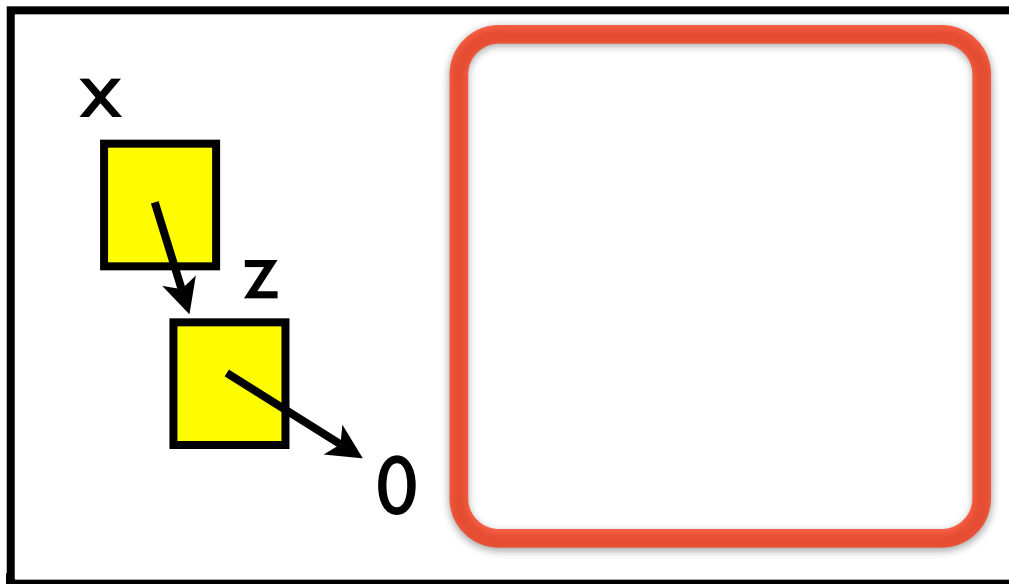
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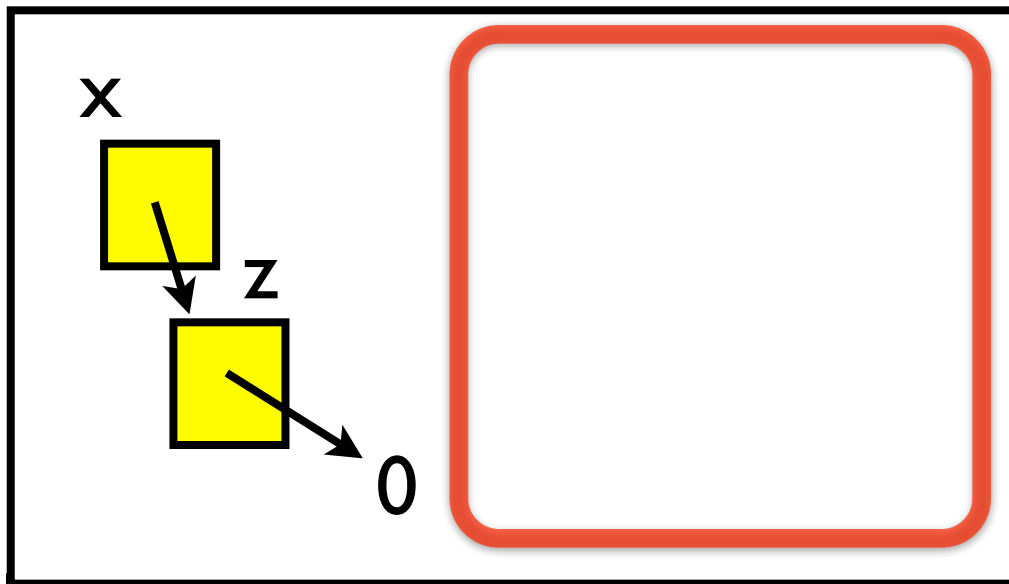
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Pre	Post
$\text{lsne } y \ y' * \text{lsne } y' \ 0$	emp

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- E.g.
 - $\text{lsne } x \ z * z \mapsto 0 * \text{lsne } y \ y' * \text{lsne } y' \ 0, \quad \text{dispose}(y)$

Symbolic Heaps with Possibly Empty List-seg Ispe

$e ::= x \mid x' \mid \text{nil}$

$\Pi ::= \Pi \wedge \Pi \mid e=e \mid e \neq e \mid \text{true}$

$\Sigma ::= \Sigma * \Sigma \mid \text{emp} \mid (e \mapsto e) \mid \text{lsne } e \ e \mid \text{lspe } e \ e \mid \text{true}$

$q ::= \Pi \wedge \Sigma$

Symbolic Heaps with Possibly Empty List-seg Ispe

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Optimization 2: Partial Join Operator

- $pjoin : SH \times SH \rightarrow SH$

- Overapproximation :

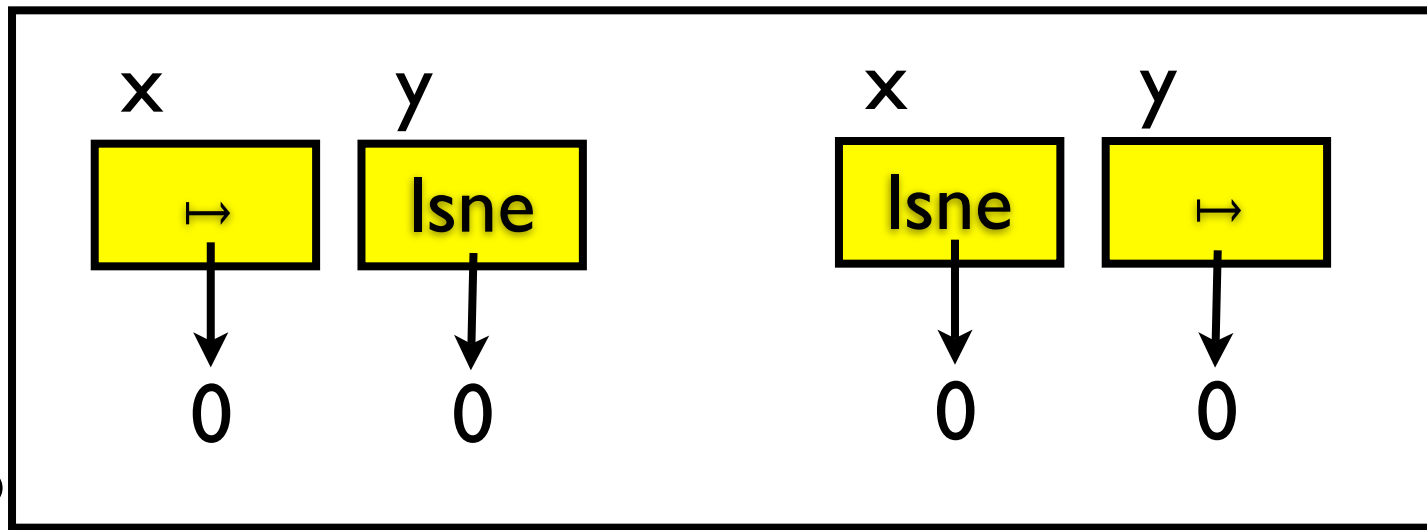
If $pjoin(q, q') = q''$, then $(q \vee q') \Rightarrow q''$

- E.g.

$$pjoin(x \mapsto 0 * \text{Isne } y \ 0, \text{Isne } x \ 0 * y \mapsto 0) = \text{Isne } x \ 0 * \text{Isne } y \ 0$$

$$pjoin(x \mapsto y * \text{Isne } y \ 0, \text{Isne } x \ 0 * y \mapsto 0) = \text{undefined.}$$

$$pjoin(y = 0 \wedge \text{Isne } x \ y, x \mapsto y * \text{Isne } y \ 0) = \text{Isne } x \ y * \text{Ispe } y \ 0$$



Operator

- Overapproximation :

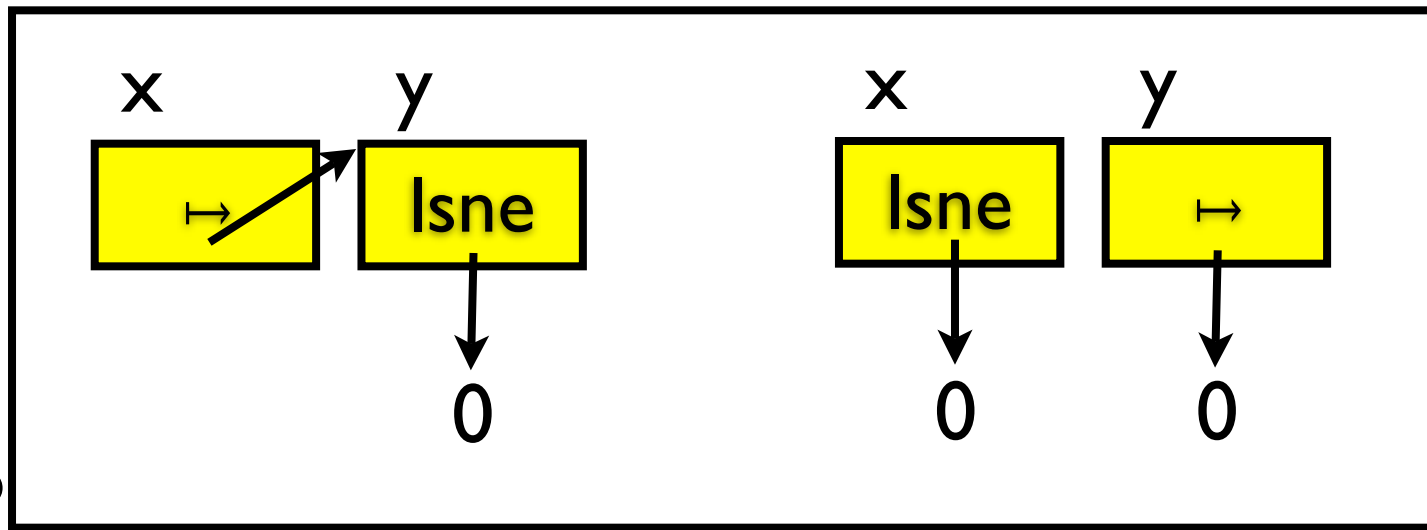
If $p \text{ join}(q, q') = q''$, then $(q \vee q') \Rightarrow q''$

- E.g.

$$p \text{ join}(x \mapsto 0 * !sne \ y \ 0, !sne \ x \ 0 * y \mapsto 0) = !sne \ x \ 0 * !sne \ y \ 0$$

$$p \text{ join}(x \mapsto y * !sne \ y \ 0, !sne \ x \ 0 * y \mapsto 0) = \text{undefined.}$$

$$p \text{ join}(y = 0 \wedge !sne \ x \ y, x \mapsto y * !sne \ y \ 0) = !sne \ x \ y * !spe \ y \ 0$$



Operator

- Overapproximation :

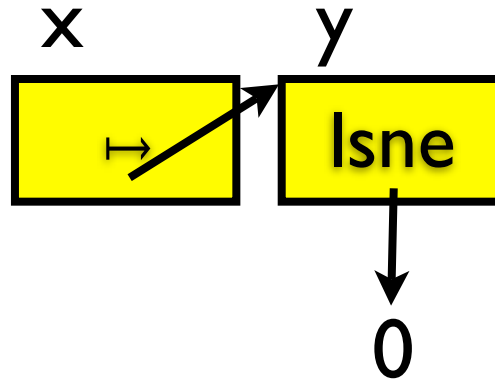
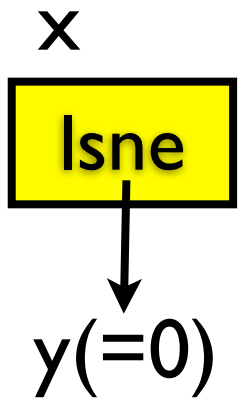
If $p \text{ join}(q, q') = q''$, then $(q \vee q') \Rightarrow q''$

- E.g.

$p \text{ join}(x \mapsto 0 * \text{Isne } y \ 0, \text{Isne } x \ 0 * y \mapsto 0) = \text{Isne } x \ 0 * \text{Isne } y \ 0$

$p \text{ join}(x \mapsto y * \text{Isne } y \ 0, \text{Isne } x \ 0 * y \mapsto 0) = \text{undefined.}$

$p \text{ join}(y = 0 \wedge \text{Isne } x \ y, x \mapsto y * \text{Isne } y \ 0) = \text{Isne } x \ y * \text{Ispe } y \ 0$



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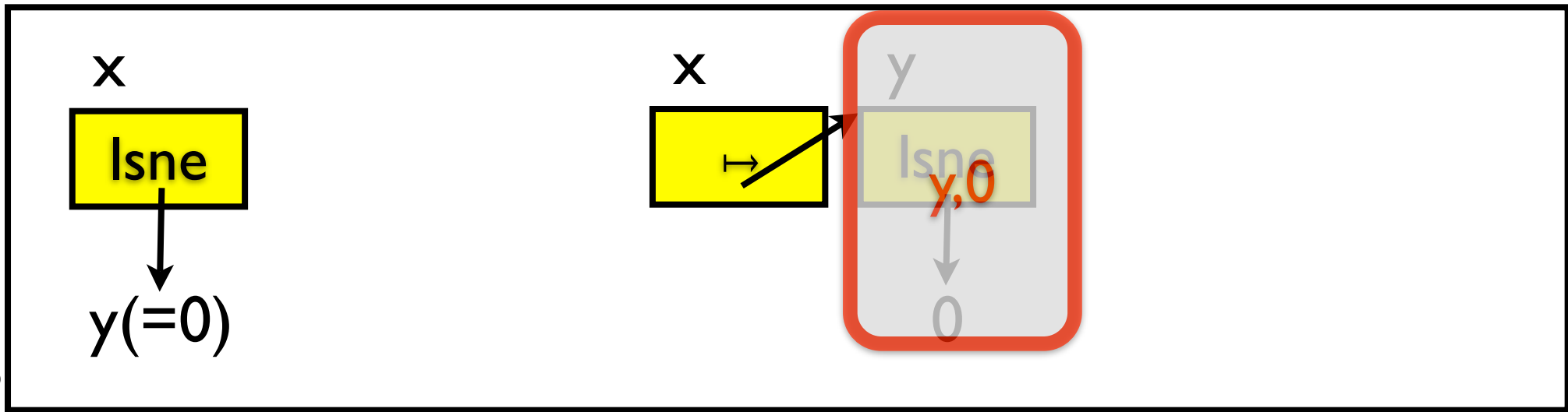
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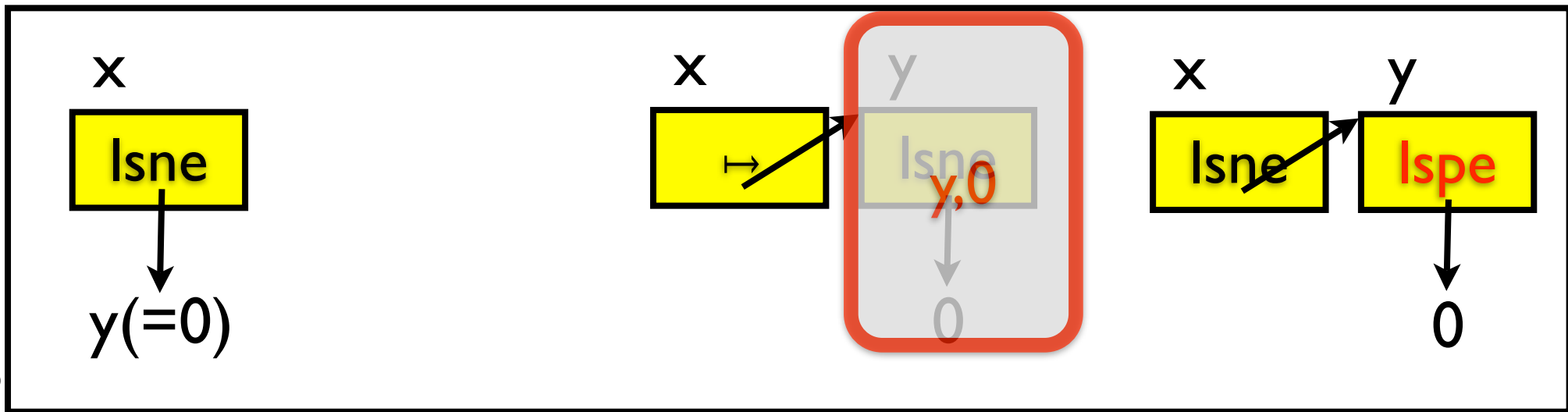
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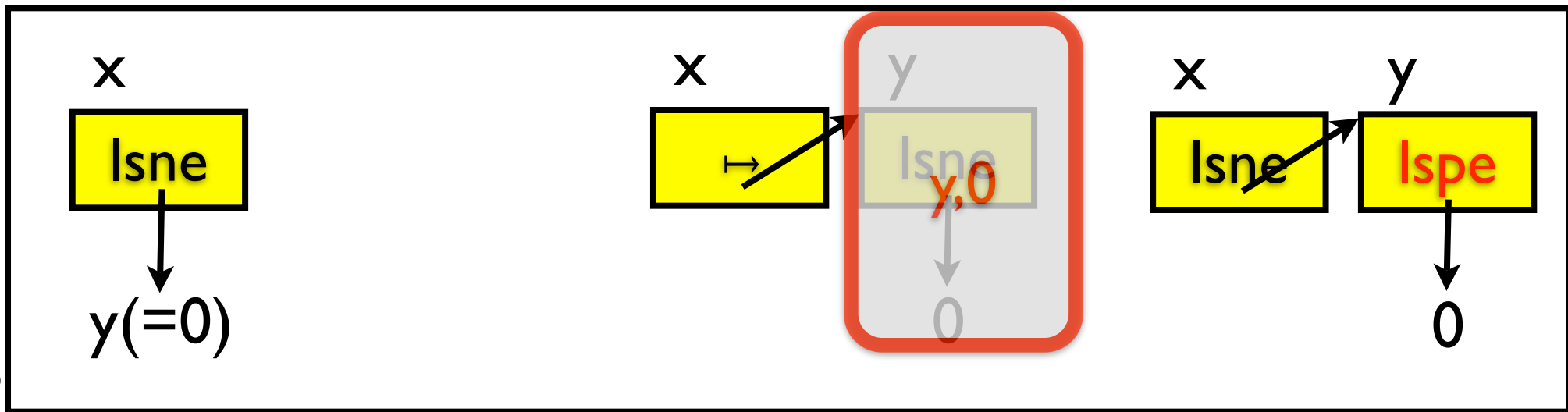
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1. Most tricky part.

2. Discovers, on the fly, which nodes to forget.

$$pjoin(x \mapsto 0 * lsne y 0, lsne x 0 * y \mapsto 0) = lsne x 0 * lsne y 0$$

$$pjoin(x \mapsto y * lsne y 0, lsne x 0 * y \mapsto 0) = \text{undefined.}$$

$$pjoin(y=0 \wedge lsne x y, x \mapsto y * lsne y 0) = lsne x y * lspe y 0$$

emp

L create() {...}

L append(L a,L b) {...}

emp

void main() {

x=create();

y=create();

x=append(x,y);

}

L create() {...}

emp

ret=0 \wedge emp,

ret \mapsto 0,

!sne ret 0

L append(L a,L b) {...}

emp

```
void main() {  
  x=create();  
  y=create();  
  x=append(x,y);  
}
```

emp

lspe ret 0

L create() {...}

L append(L a,L b) {...}

emp

```
void main() {  
  x=create();  
  y=create();  
  x=append(x,y);  
}
```

emp

lspe ret 0

L create() {...}

L append(L a,L b) {...}

emp

void main() {

x=create(); lspe x 0

y=create();

x=append(x,y);

}

emp

lspe ret 0

```
L create() {...}
```

```
L append(L a,L b) {...}
```

emp

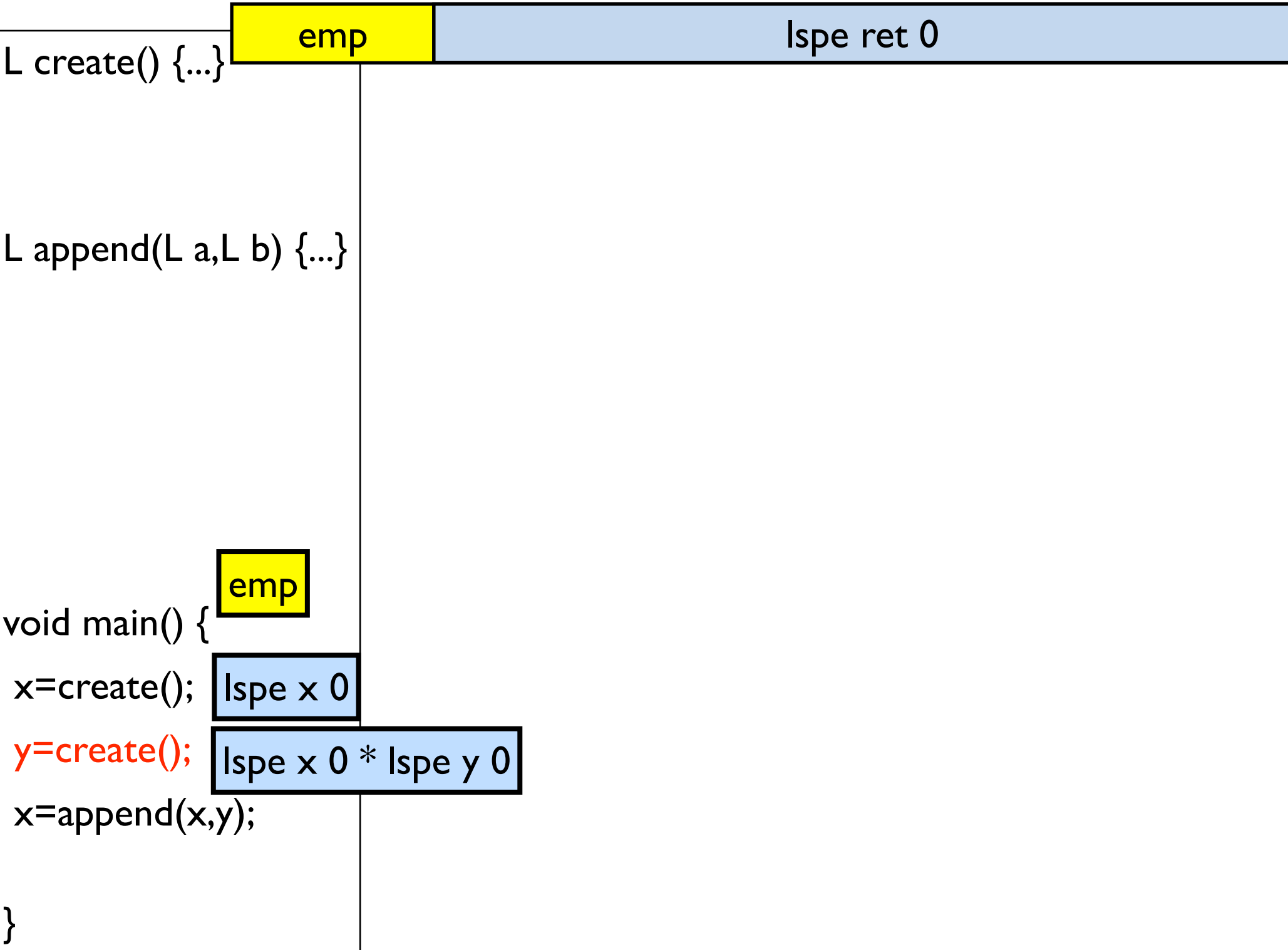
```
void main() {
```

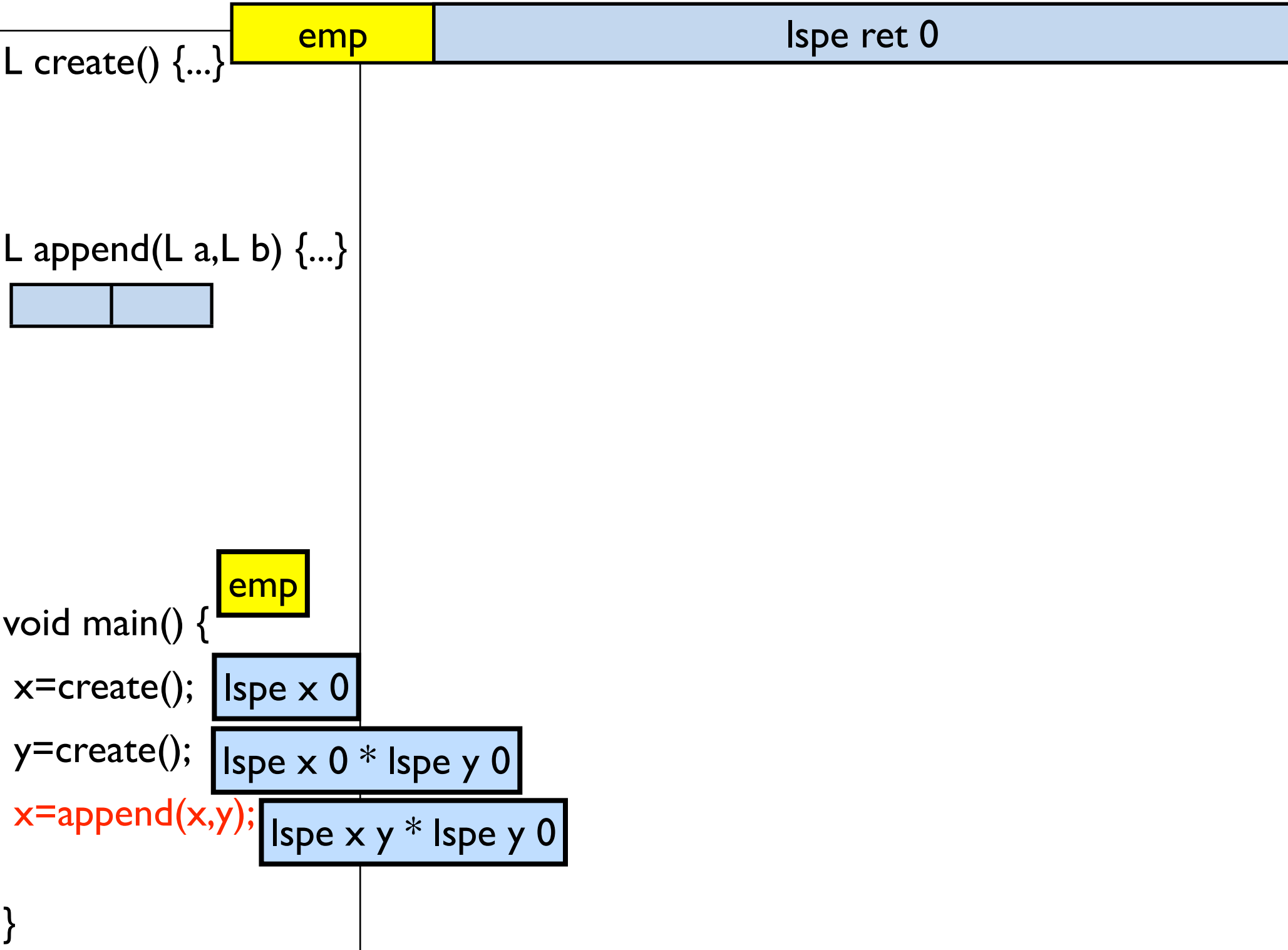
```
  x=create(); lspe x 0
```

```
  y=create();
```

```
  x=append(x,y);
```

```
}
```





L create() {...}

emp

lspe ret 0

~~4 entries, 12 results~~

L append(L a,L b) {...}



~~9 entries, 12 results~~

All tables have
one entry, one result

emp

void main() {

x=create(); lspe x 0

y=create(); lspe x 0 * lspe y 0

x=append(x,y); lspe x y * lspe y 0

}

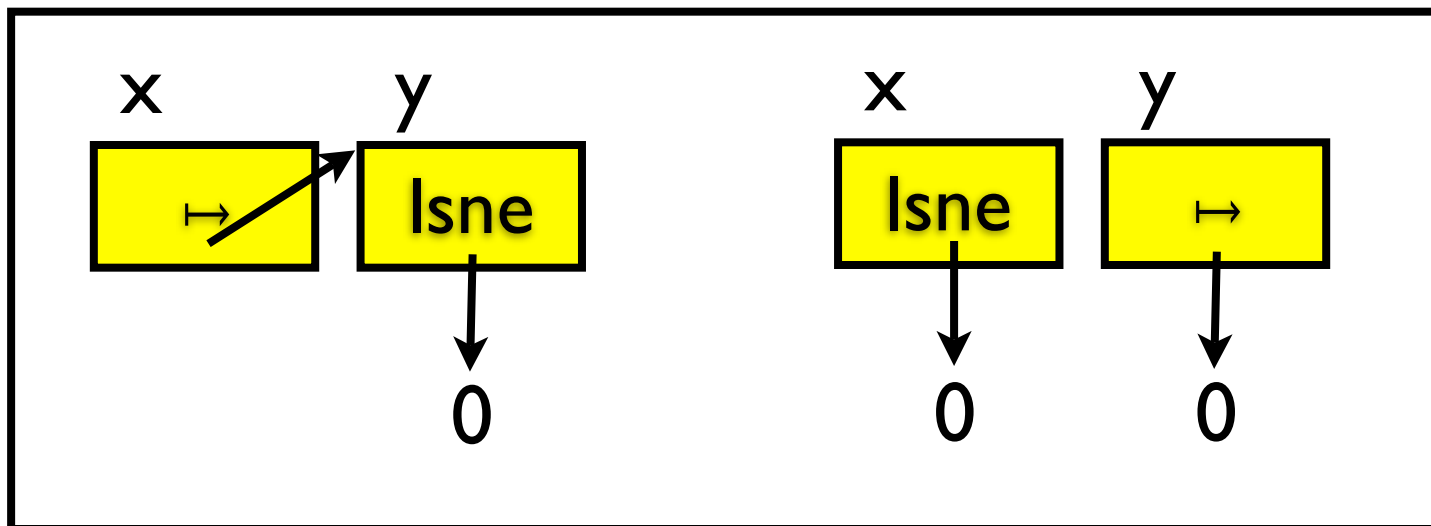
Why Partial Join?

Prevent the analysis from misusing existential variables and losing crucial shape information.

$$\{ x \mapsto y * \text{lsne } y 0 \vee \text{lsne } x 0 * y \mapsto 0 \}$$

`free_list(x)`

$$\{ \text{emp} \vee y \mapsto 0 \}$$

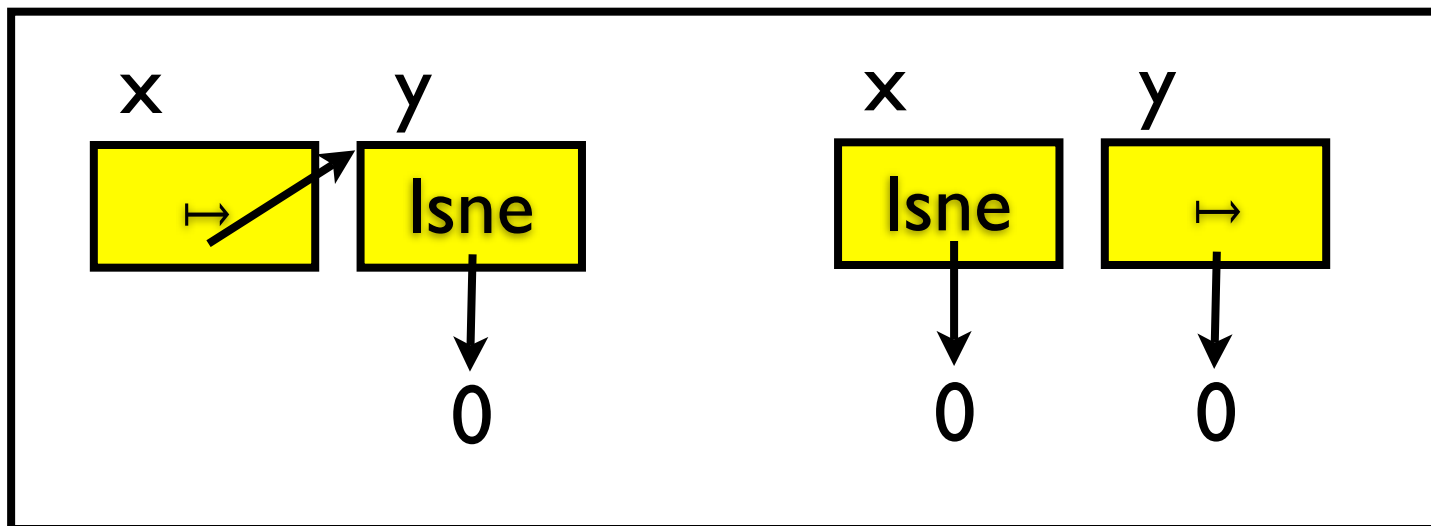


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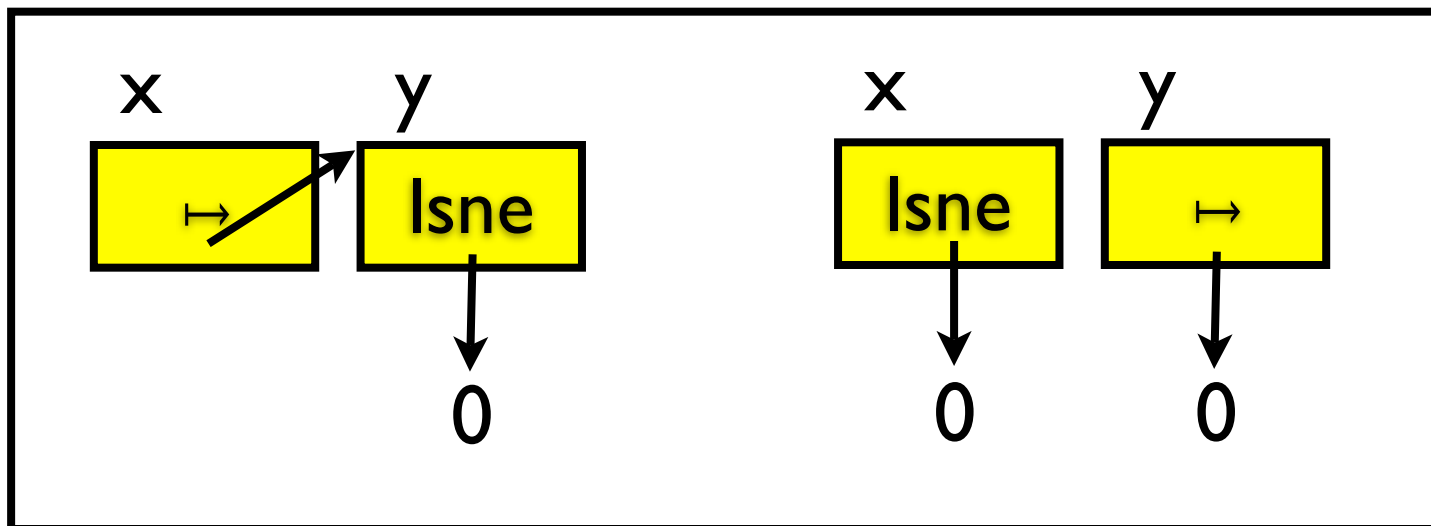


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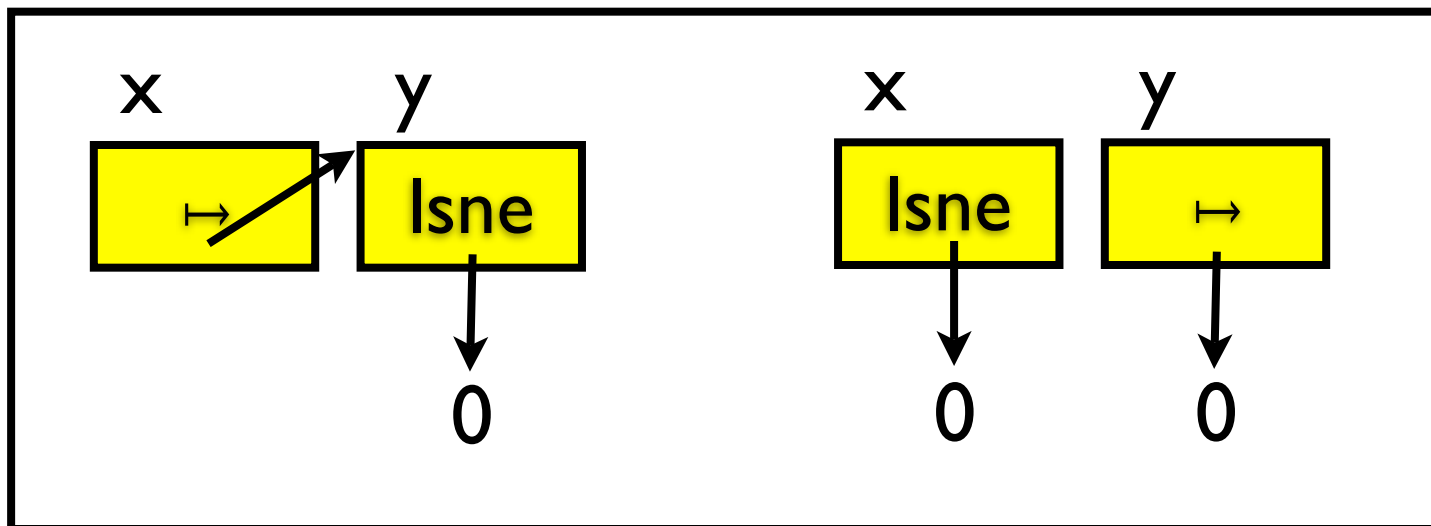


Prevent the analysis from misusing existential variables and losing crucial shape information.

```

{   lsne x x' * lsne y 0   }
  free_list(x)
{   }

```



Prevent the analysis from misusing existential variables and losing crucial shape information.

```

{   lsne x x' * lsne y 0   }
  free_list(x)
{   ???   }

```

Why Partial Join?

Keep important co-relation between value and shape.

```
{ x ↦ 0 }
```

```
y = t1394Diag_IoControl(x);
```

```
{ y = 0 * emp ∨ y != 0 * x ↦ 0 }
```

```
if (y != 0) free(x);
```

```
{ emp }
```

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```
{ ???? }
```

Ingredients in Solution

- Partial join for separation domain (CAV08).
- Composite adaptive analysis (CAV07).
- Localization (CSL01, Rinetzky+POPL05, Gotsman+SAS06).
- Partial concretization (Sagiv+TOPLAS98).
- RHS (Reps+POPL95).
- Separation logic.

Open Question, Summer 2006

Can we automatically prove the pointer safety of programs $\geq 10K$ in separation logic?