

### C340 Concurrency: Starvation and Liveness

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Reader/Writer problem

#### Starvation

- Dining Philosophers Problem
- Deadlocks
- Liveness Analysis using LTS



**Reader / Writer Problem** 

Monitors and Java's synchronize statement guarantee mutual access to objects / methods

Often it is ok for multiple readers to access the object concurrently

Properties required:





### **Read/Write Monitor**

```
class ReadWrite {
private protected int readers = 0;
private protected boolean writing = false;
 // Invariant: (readers>=0 and !writing) or
 // (readers==0 and writing)
 synchronized public void acquireRead() {
 while (writing) {... wait(); ...} ++readers;
 synchronized public void releaseRead() {
    --readers; if(readers==0) notify();
 synchronized public void acquireWrite() {
 while (readers>0||writing) {... wait(); ...}
 writing = true;
 synchronized public void releaseWrite() {
 writing = false; notifyAll();
                                   Starvation
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```



- NotifyAll awakes both readers and writers
- Program relies on Java having a fair scheduling strategy
- When readers continually read resource: Writer never gets chance to write. This is an example of starvation.
  - Solution: Avoid writer starvation by making readers defer if there is a writer waiting



# Read/Write Monitor (Version 2)

```
class ReadWrite {
... // as before
private int waitingW = 0;// # waiting Writers
 synchronized public void acquireRead() {
 while (writing || waitingW>0) {... wait(); ... }
  ++readers;
synchronized public void releaseRead() {... }
 synchronized public void acquireWrite() {
 while (readers>0 || writing) {
  ++waitingW; ... try{ wait(); ... --waitingW; }
 writing = true;
synchronized public void releaseWrite() {... }
                     Demo: Reader/Writer v2
```

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**Reader Starvation** 

If there is always a waiting writer: Readers starve Solution: Alternating preference between readers and writers To do so: Another boolean attribute readersturn in Monitor that indicates whose turn it is readersturn is set by releaseWrite() and cleared by releaseRead()



# Read/Write Monitor (Version 3)

```
class ReadWrite {
... // as before
private boolean readersturn = false;
 synchronized public void acquireRead() {
 while(writing ||(waitingW>0 && !readersturn))
     { ... wait(); ... }
  ++readers;
 synchronized public void releaseRead() {
  --readers; readersturn=false;
  if(readers==0) notifyAll();
 synchronized public void acquireWrite() {... }
 synchronized public void releaseWrite() {
 writing=false; readersturn=true; notifyAll();
                   Demo: Reader/Writer v3
                                                 8
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```



### Deadlocks

- Process is in a <u>deadlock</u> if it is blocked waiting for a condition that will never become true
- Process is in a <u>livelock</u> if it is spinning while waiting for a condition that will never become true (busy wait deadlock)
- Both happen if concurrent processes and threads are mutually waiting for each other

#### Example: Dining philosophers

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# **Dining Philosopher Problem**

- 5 Philosophers sit around table
  - They think or eat
- Eat with 2 chopsticks
- Only 5 chopsticks available
- Each philosopher only uses sticks to her left and right





## FSP Model of Dining Philosophers





# **Dining Philosophers in Java**

```
class Philosopher extends Thread {
  int identity;
  Chopstick left; Chopstick right;
  Philosopher(Chopstick left, Chopstick right) {
    this.left = left; this.right = right;
  public void run() {
  while (true) {
    try {
     sleep(...);
                               // thinking
     right.get(); left.get(); // hungry
     sleep(...) ;
                              // eating
     right.put(); left.put();
    { catch (InterruptedException e) { }
```



## **Chopstick Monitor**

```
class Chopstick {
 boolean taken=false;
 synchronized void put() {
  taken=false;
  notify();
 synchronized void get() throws
                     InterruptedException
  while (taken) wait();
     taken=true;
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```



### **Applet for Diners**

#### for (int i =0; i<N; ++I)</pre> // create Chopsticks stick[i] = new Chopstick(); for (int i =0; i<N; ++i){</pre> // create Philosophers phil[i]=new Philosopher( stick[(i-1+N)%N],stick[i]); phil[i].start(); } Demo: Diners



- If each philosopher has acquired her left chopstick the threads are mutually waiting for each other
- Potential for deadlock exists independent of thinking and eating times
- Only probability is increased if these times become shorter



- We can use LTS for deadlock analysis
- A <u>dead state</u> in the composed LTS is one that does not have outgoing transitions
- Are these dead states reachable?
- Use of reachability analysis
- Traces to dead states helps understanding the causes of a deadlock





Deadlock Avoidance

Deadlock in dining philosophers can be avoided if one philosopher picks up sticks in reverse order (right before left).

Demo: Deadlock free Diners

What is the problem with this solution?
Are there other solutions?
Deadlock can also be avoided there is always one philosopher who thinks



### Summary

#### Reader / Writer Problem

#### Starvation

- Avoidance of Starvation
  - Dining Philosophers Problem
  - **Deadlocks and Livelocks**
  - Deadlock Avoidance
- Next Session: Safety