An introduction to separation logic

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Introduction

Verification of imperative programs is classically based on Hoare triples:

\[ \{P\} \ C \ \{Q\} \]

where \( C \) is a program and \( P, Q \) are assertions in some logical language.
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These are read, roughly speaking, as

for any state \( \sigma \) satisfying \( P \), if \( C \) transforms state \( \sigma \) to \( \sigma' \), then \( \sigma' \) satisfies \( Q \).
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\[ \text{for any state } \sigma \text{ satisfying } P \text{, if } C \text{ transforms state } \sigma \text{ to } \sigma' \text{, then } \sigma' \text{ satisfies } Q. \]

(with some wriggle room allowing us to deal with faulting or non-termination in various ways.)
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3. a formal interpretation of Hoare triples, together with (sound) proof rules for manipulating them.
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3. a formal interpretation of Hoare triples, together with (sound) proof rules for manipulating them.

We’ll look at these informally first, then introduce a little more formal detail.
Programs, informally

We consider a standard \textbf{while} language with \textbf{pointers}, memory (de)\textbf{allocation} and recursive \textbf{procedures}.
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```plaintext
deltree(*x) {
    if x=nil then return;
    else {
        l,r := x.left,x.right;
        deltreen(l);
        deltreen(r);
        free(x);
    }
}
```

Assertions, informally

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E.g., binary trees with root pointer $x$ can be defined by:

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x = \text{nil} : \text{emp} & \Rightarrow \text{tree}(x) \\
x \neq \text{nil} : x & \mapsto (y, z) \ast \text{tree}(y) \ast \text{tree}(z) \Rightarrow \text{tree}(x)
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where

- $\text{emp}$ denotes the empty heap;
- $x \mapsto (y, z)$ denotes a single pointer to a pair of data cells;
- $\ast$ means “and separately in memory”.

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An example proof

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\}
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\end{align*}
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An example proof

\{\text{tree}(x)\}
\text{deltree}(\ast x) \{
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  \text{else } \{x \mapsto (y, z) \ast \text{tree}(y) \ast \text{tree}(z)\}
  l, r := x.\text{left}, x.\text{right};
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\begin{verbatim}
{tree(x)}
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        {x \mapsto (l, r) * \text{emp} * \text{tree}(r)}
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Consider the following step in the previous example:

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\[
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Consider the following step in the previous example:

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\text{deltree}(l)
\]
\[
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\]

Implicitly, this relies on a **framing property**, namely:

\[
\{\text{tree}(l)\} \text{deltree}(l) \{\text{emp}\}
\]
\[
\{x \mapsto (l, r) \ast \text{tree}(l) \ast \text{tree}(r)\} \text{deltree}(l) \{x \mapsto (l, r) \ast \text{emp} \ast \text{tree}(r)\}
\]
Classical failure of frame rule

The so-called frame rule,

\[
\begin{align*}
\{P\} & \ C \ \{Q\} \\
\{F \land P\} & \ C \ \{F \land Q\}
\end{align*}
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is well known to fail in standard Hoare logic.
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\begin{array}{c}
\{x = 0\} \ x := 2 \ \{x = 2\} \\
\hline
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is not valid (because \(y\) could alias \(x\)).
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As we’ll see, using the “separating conjunction” \( \ast \) instead of \( \land \) will however give us a valid frame rule.
Heap memory model

- We assume an infinite set $\text{Val}$ of values of which an infinite subset $\text{Loc} \subset \text{Val}$ are allocable locations; nil is a non-allocable value.
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• Heaps map finitely many locations to values, $h : \text{Loc} \rightarrow_{\text{fin}} \text{Val}$. We write $e$ for the empty heap (undefined on all locations).
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• Heap composition $h_1 \circ h_2$ is defined to be $h_1 \cup h_2$ if their domains are non-overlapping, and undefined otherwise.
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• A state is simply a stack paired with a heap, \((s, h)\).
Program semantics

- A **configuration** is given by \((C, s, h)\), where \(C\) is a program, and \((s, h)\) a (stack-heap) state.
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• fault is a special configuration used to catch memory errors.

• The small-step semantics of programs is then given by a relation \(\rightsquigarrow\) between configurations:

\[(C, s, h) \rightsquigarrow (C', s', h')\]
Semantics of assignment and (de)allocation

\[
(x := E, s, h) \leadsto (s[x \mapsto \sem E s], h)
\]

\[
\sem E s \in \text{dom}(h)
\]

\[
(x := E.f, s, h) \leadsto (s[x \mapsto h(\sem E s).f], h)
\]

\[
\sem E s \in \text{dom}(h)
\]

\[
(E.f := E', s, h) \leadsto (s, h[\sem E s.f \mapsto \sem E' s])
\]

\[
\ell \in \text{Loc} \setminus \text{dom}(h) \quad v \in \text{Val}
\]

\[
(E := \text{new}(), s, h) \leadsto (s[x \mapsto \ell], h[\ell \mapsto v])
\]

\[
\sem E s = \ell \in \text{dom}(h)
\]

\[
(free(E), s, h) \leadsto (s, (h \upharpoonright (\text{dom}(h) \setminus \{\ell\}))
\]

\[
C \equiv x := E.f \mid E.f := E' \mid free(E) \mid \sem E s \notin \text{dom}(h)
\]

\[
(C, s, h) \leadsto \text{fault}
\]

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Symbolic-heap assertions

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- Pure formulas $\pi$ and spatial formulas $F$ are given by:

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\pi ::= t = t \mid t \neq t \\
F ::= \text{emp} \mid x \mapsto t \mid P\mathbf{t} \mid F \ast F
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(where $P$ a predicate symbol, $\mathbf{t}$ a tuple of terms).
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- A symbolic heap is $\exists x. \Pi : F$, for $\Pi$ a set of pure formulas.
- The predicate symbols might come from a hard-coded set, or might be user-defined.
Semantics of assertions

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- $s, h \models \phi \, t_1 = (\neq) t_2 \iff s(t_1) = (\neq)s(t_2)$
- $s, h \models \phi \, \text{emp} \iff h = e$
- $s, h \models \phi \, x \mapsto t \iff \text{dom}(h) = \{s(x)\}$ and $h(s(x)) = s(t)$
- $s, h \models \phi \, P_\mathbf{t} \iff (s(t), h) \in [P]$
- $s, h \models \phi F_1 \times F_2 \iff \exists h_1, h_2. \ h = h_1 \circ h_2$ and $s, h_1 \models \phi F_1$
  and $s, h_2 \models \phi F_2$
- $s, h \models \phi \exists z. \Pi : F \iff \exists v \in \text{Val}^{\lvert z \rvert}. \ s[z \mapsto v], h \models \phi \pi$ for all
  $\pi \in \Pi$ and $s[z \mapsto v], h \models \phi F$

The semantics $[P]$ of inductive predicate $P$ has a standard construction (but outside the scope of this talk).
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Interpretation of Hoare triples

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\( \{ P \} C \{ Q \} \) is valid if, whenever \( s, h \models P \),

1. \((C, s, h) \not\to^{*} fault\) (i.e. is memory-safe), and
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2. if \((C, s, h) \sim^* (\epsilon, s, h)\), then \(s, h \models Q\).

If we are interested in total correctness, simply replace the memory-safety condition above by (safe) termination: everything still works!
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If we are interested in total correctness, simply replace the memory-safety condition above by (safe) termination: everything still works!
Axioms and proof rules for triples

\[
\begin{align*}
\{\text{emp}\} \ x := E \ \{x = E[x'/x] : \text{emp}\} & \quad \{E.f \mapsto \_\} \ E.f := E' \ \{E.f \mapsto E'\} \\
\{E.f \mapsto t\} \ x := E.f \ \{x = t[x'/x] : E.f \mapsto t[x'/x]\} & \\
\{\text{emp}\} \ x := \text{new()} \ \{x \mapsto x'\} & \quad \{E \mapsto \_\} \ \text{free}(E) \ \{\text{emp}\} \\
\{P\} \ C_1 \ \{R\} \quad \{R\} \ C_2 \ \{Q\} & \quad \{B : P\} \ C_1 \ \{Q\} \quad \{\neg B : P\} \ C_2 \ \{Q\} \\
\{P\} \ C_1; C_2 \ \{Q\} & \quad \{P\} \ \text{if } B \ \text{then } C_1 \ \text{else } C_2 \ \{Q\}
\end{align*}
\]

(Note that \(E.f \mapsto E'\) is a shorthand for \(E \mapsto (\ldots, E', \ldots)\) where \(E'\) occurs at the \(f\)th position in the tuple.)
The general frame rule of separation logic can be stated as follows:

\[
\frac{\{P\} \ C \ \{Q\}}{\{F \times P\} \ C \ \{F \times Q\}}
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subject to the obvious sanity condition: \( C \) does not modify any variable mentioned in the “frame” \( F \).
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subject to the obvious sanity condition: \(C\) does not modify any variable mentioned in the “frame” \(F\).

This rule is exactly what is needed to carry out proofs like the one we saw before for \texttt{deltree}.
Soundness of frame rule

**Soundness** of the frame rule depends on the following two facts about the programming language:

**Lemma (Safety monotonicity)**
If \((C,s,h) \not\leadsto \ast\text{fault}\) then \((C,s,h') \not\leadsto \ast\text{fault}\) (for any \(h'\) such that \(h \circ h'\) is defined).

**Lemma (Frame property)**
Suppose \((C,s,h_1 \circ h_2) \leadsto \ast\langle s,h_1 \rangle\), and that \((C,s,h_1) \not\leadsto \ast\text{fault}\). Then there exists \(h'\) such that \((C,s,h_1) \leadsto \ast\langle s,h_2\rangle\), and, moreover, \(h' = h \circ h_2\).

Together, these lemmas imply the locality of all commands.

N.B.: this is an operational fact about the programming language, and nothing at all to do with logic!
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**Lemma (Safety monotonicity)**

If $(C, s, h) \not fault$ then $(C, s, h \circ h') \not fault$ (for any $h'$ such that $h \circ h'$ is defined).

**Lemma (Frame property)**

Suppose $(C, s, h_1 \circ h_2) \leadsto* \langle s, h \rangle$, and that $(C, s, h_1) \not fault$. Then there exists $h'$ such that $(C, s, h_1) \leadsto* \langle s, h' \rangle$, and, moreover, $h = h' \circ h_2$. 
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Then there exists \(h'\) such that \((C, s, h_1) \leadsto^* \langle s, h' \rangle\), and, moreover, \(h = h' \circ h_2\).

Together, these lemmas imply the locality of all commands. N.B.: this is an operational fact about the programming language, and nothing at all to do with logic!
Closing remarks

- What we call separation logic is really a combination of
  - programming language,
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  - and rules for Hoare triples.
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  - programming language,
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  - and rules for Hoare triples.

- The power of separation logic comes from compositionality: proofs of sub-programs can be combined into proofs of whole programs.

- Compositionality depends on the frame rule:

\[
\begin{align*}
\{P\} C \{Q\} \\
\{F \ast P\} C \{F \ast Q\}
\end{align*}
\]
Closing remarks

- What we call separation logic is really a combination of
  - programming language,
  - assertion language
  - and rules for Hoare triples.

- The power of separation logic comes from compositionality: proofs of sub-programs can be combined into proofs of whole programs.

- Compositionality depends on the frame rule:

  \[
  \begin{array}{c}
  \{P\} C \{Q\} \\
  \hline
  \{F \ast P\} C \{F \ast Q\}
  \end{array}
  \]

- And the soundness of the frame rule is essentially a reflection of the locality of commands.
Further reading

S. Ishtiaq and P. O’Hearn.
BI as an assertion language for mutable data structures.
(Winner of Most Influential POPL Paper 2001 award.)

J.C. Reynolds.
Separation logic: A logic for shared mutable data structures.

H. Yang and P. O’Hearn.
A semantic basis for local reasoning.

Compositional shape analysis by means of bi-abduction.