

# **Network Security: Internet Worms and Firewalls**

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# Outline

- Internet worms
  - Self-propagating, possibly malicious code spread over Internet
- Firewalls:
  - Simple, perimeter-based security

# What's a Worm?

- Vast numbers of Internet-attached hosts run **vulnerable server software**
- Worm: **self-replicating code**, containing
  - **Exploit** for widely used, vulnerable server software
  - **Payload**: code that executes after exploit succeeds
- Payload connects to other Internet hosts, **sends copy of {exploit, payload} to each...**
- Unlike virus, **spread not human-mediated**

# What's in the Payload?

- Could be anything...arbitrary code execution allowed by many exploits
- Install login facility for attacker, to allow use at will in botnet
  - Botnets used widely today to launch DDoS attacks, send spam
  - Market in botnets exists today (3-10 US cents/host/week for spam proxy in 2005 [Paxson])
- Send sensitive files to attacker
- Destroy or corrupt data
- Enormous possibility for harm, in financial, privacy, and inconvenience terms

# Code-RedI Worm

- June 18<sup>th</sup>, 2001: eEye releases description of **buffer overflow vulnerability in Microsoft IIS (web server)**
- June 26<sup>th</sup>, 2001: **Microsoft releases patch**
- July 12<sup>th</sup>, 2001: **Code-RedI worm released** (i.e., first sent to vulnerable host)
- Estimated number hosts infected: **360,000**
- Estimated damages: **\$2.6 billion** from loss of service availability, downtime, cleanup...

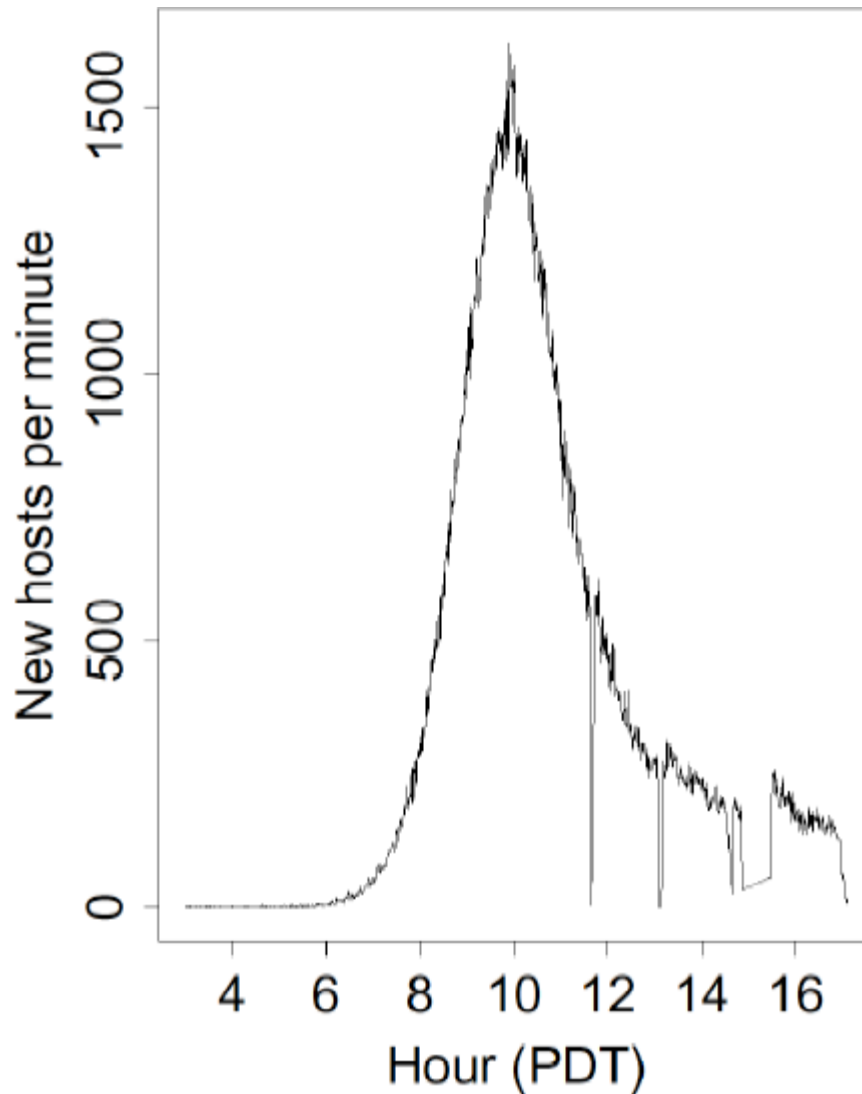
# Code-RedI Behavior

- Payload: defaces web site
  - If language == English
    - HELLO! Welcome to <http://www.worm.com>!  
Hacked By Chinese!
- 1<sup>st</sup> – 19<sup>th</sup> of every month: spread
  - Connect to random 32-bit IP address, send copy of self (exploit+payload)
- 20<sup>th</sup> through end of every month:
  - Flood traffic to 198.137.240.91  
([www.whitehouse.gov](http://www.whitehouse.gov))
- Bug: fixed seed for random number generator
  - All hosts generate same sequence of IPs!
  - Result: only linear growth in infected population
- Only memory-resident; vanishes on reboot

# Code-RedI v2: “Bugfix” Release

- July 19<sup>th</sup>, 2001: new variant (“v2”) released
  - Uses random seed
  - Now all infected hosts try different targets
- White House changes IP address of its server to avoid DDoS attack
  - Result: July 20<sup>th</sup>, Code-RedI v2 dies out
- 360K hosts infected in 14 hours

# Growth of Code-RedI v2



- Source:  
Vern Paxson,  
ICSI/UC Berkeley



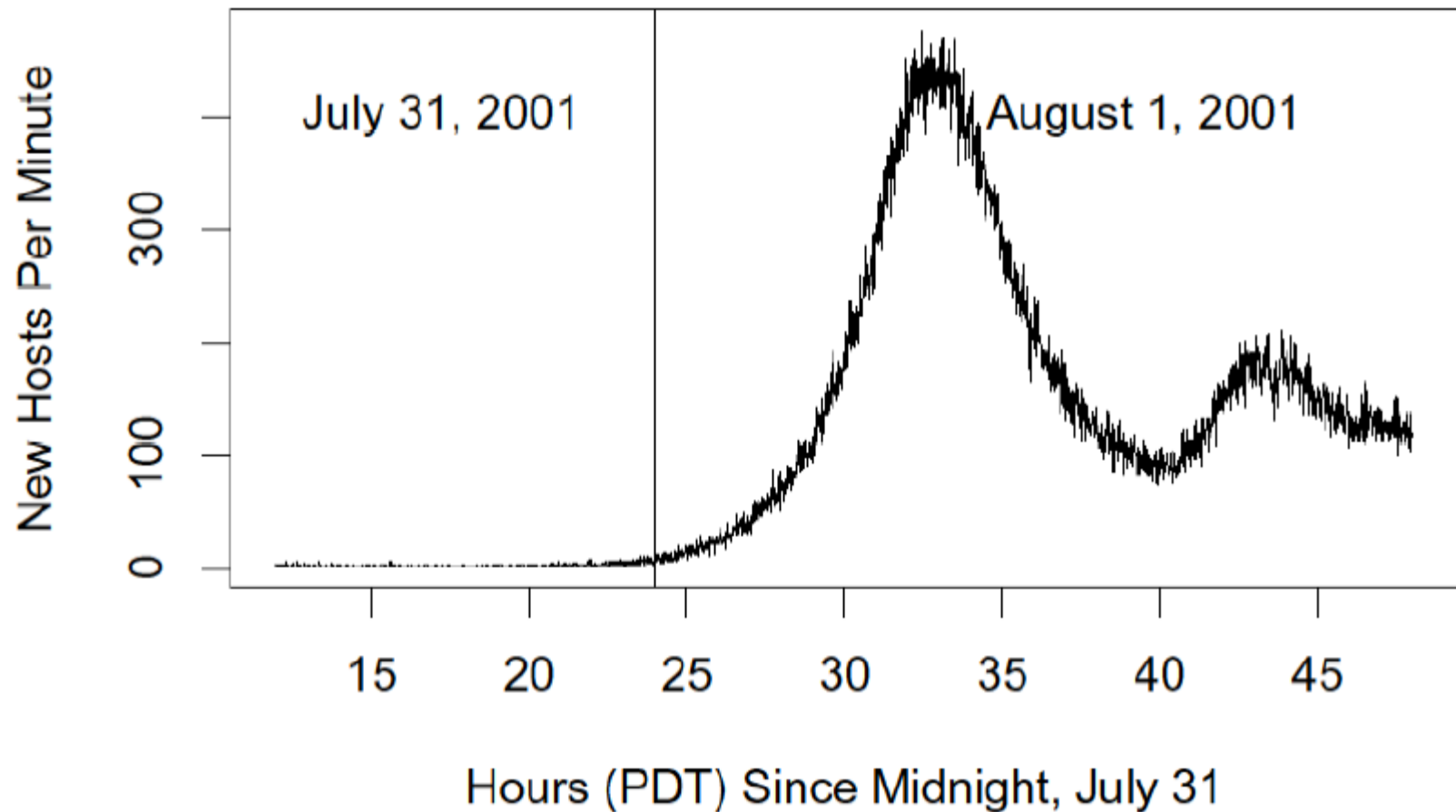
# Network Telescopes

- Monitor traffic arriving at sizeable regions of Internet address space. Reveals, e.g.,:
  - “Backscatter” (responses to randomly source-spoofed DDoS attacks)
  - Worms’ random scanning of IP addresses
  - Attackers’ random scanning for servers running particular service
- LBNL: 2 /16 networks, or  $1/32768^{\text{th}}$  of Internet address space
- UCSD/Univ. Wisconsin: 1 /8 network, or  $1/256^{\text{th}}$  of Internet address space

# Spread of Code-RedI v2

- Network telescope estimate of infected host count:
  - Count unique source IPs that attempt to connect to port 80 on non-used addresses
- Infected population over time fits logistic function
  - S-shaped curve: exponential growth at start, then slowing growth after most vulnerable nodes infected
- Worm dies just as 20<sup>th</sup> starts
  - But even one host with wrong clock can keep trying to infect others
  - On August 1<sup>st</sup>, worm begins to spread again!

# Return of Code Red Worm



- Source: Vern Paxson, ICSI/UC Berkeley

# A Competitor: Code-Red II

- Targets same IIS vulnerability; unrelated code
- Released August 4<sup>th</sup>, 2001
- Installs superuser backdoor; **persists after reboot**
- Spreads preferentially to **local addresses**:
  - 1/2 probability generates address on same /8
  - 3/8 probability generates address on same /16
  - 1/8 probability generates random non-class-D, non-loopback address
- Result: **squeezes out Code-Red I v2!**

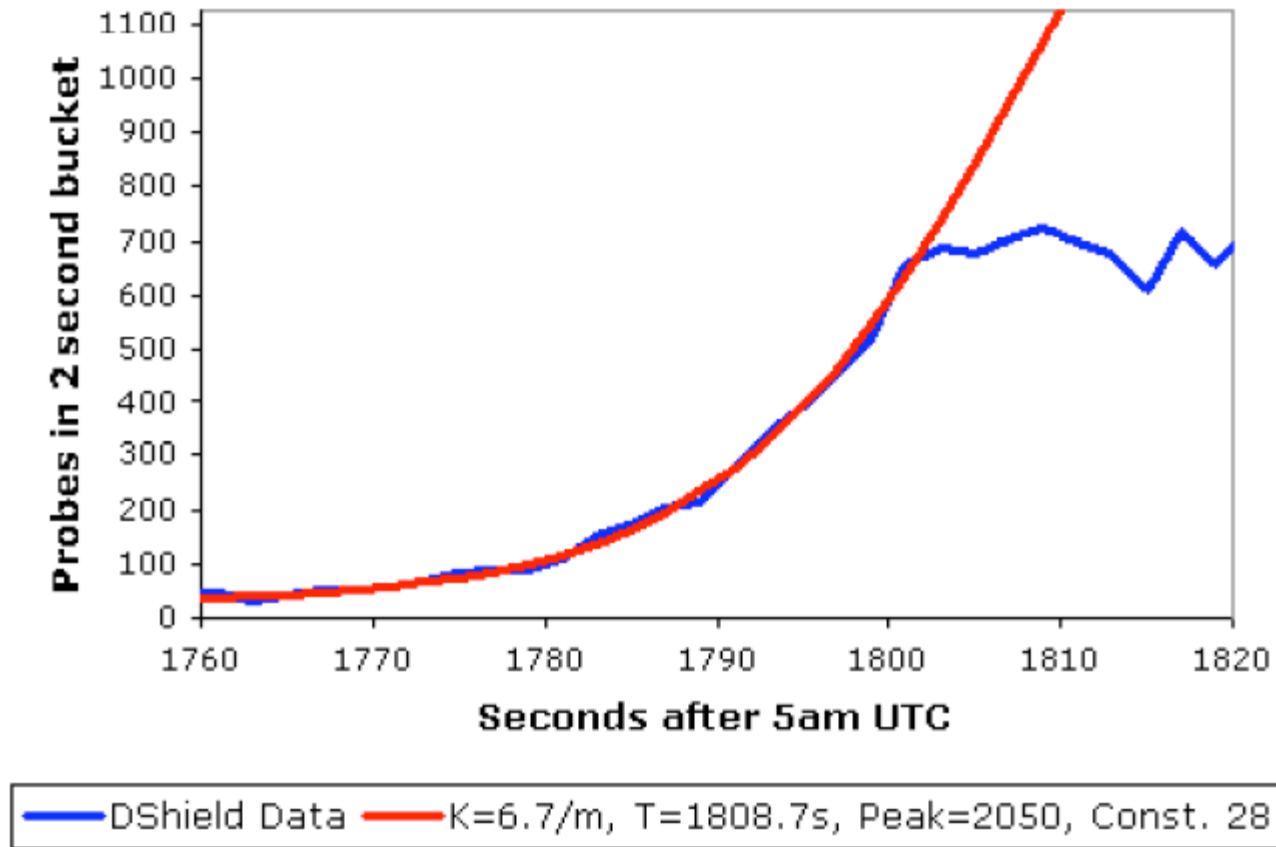
# Slammer: A Fast UDP Worm

- Exploit: buffer overflow vulnerability in Microsoft SQL Server 2000
  - Vulnerability reported in June 2002
  - Patch released July 2002
- SQL service uses **connectionless UDP** (rather than connection-oriented TCP)
- Entire worm **fit in one packet!**
  - No need to wait for RTT; send single packet, try next target address
- Slammer infected **over 75K hosts in 10 minutes**
- Growth rate **limited by Internet's capacity**

# Slammer's Behavior

- Peak address scanning rate: 55 million scans / second
  - Reached in 3 minutes
  - Beyond that point, congestion-limited
- Payload non-malicious, apart from aggressive scanning
- Outages in 911 (emergency telephone) service, Bank of America ATM network
  - Purely from traffic load; crashed some network equipment, saturated some bottleneck links

# Slammer's Growth Limited by Internet Bandwidth (!)



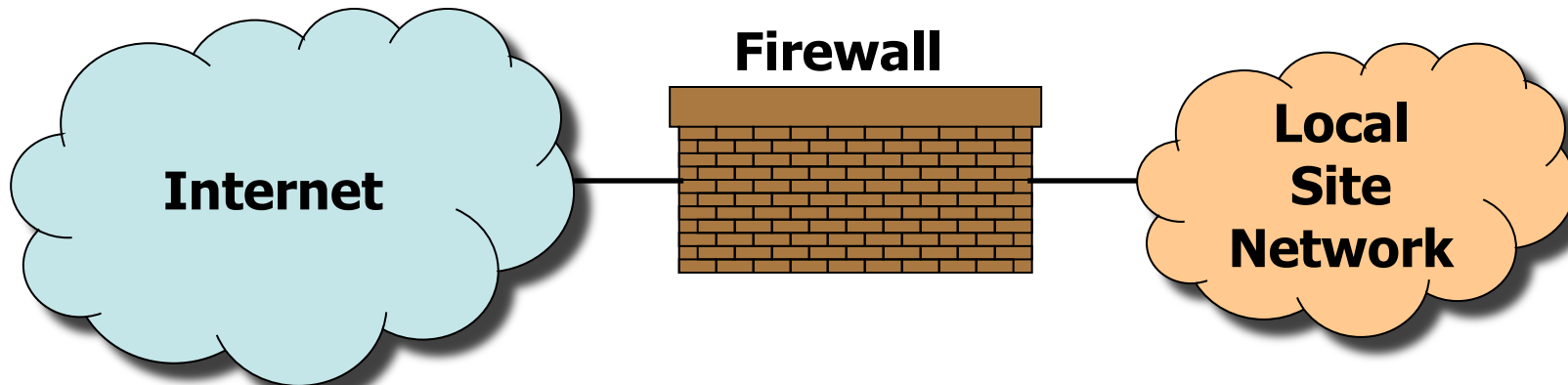
- Source: Vern Paxson, ICSI/UC Berkeley

# Worm Propagation Methods

- **Random scanning** (e.g., Code-Red, Slammer)
- **Meta-server worm**: query a service for hosts to infect (e.g., ask Google, “powered by phpbb”)
- **Topological worm**: find candidates from files on infected host’s disk (e.g., web server logs, bookmark files, email address books, ssh known hosts files, ...)
  - Very fast; stealthy—no random scanning behavior to attract attention
- **Contagion worm**: piggyback worm on application’s usual connections
  - Connection patterns appear normal!



# Firewalls: Perimeter-Based Defense



- Define **trusted perimeter** (typically boundary of own infrastructure)
- All packets between Internet and trusted perimeter flow through **firewall**
- Firewall **inspects, filters traffic** to **limit access to non-secure services** by remote, untrusted hosts

# Firewall: Physical Topology vs. Filtering Policies

- Topological placement of firewall depends on **perimeter at which defense desired**, e.g.,
  - Firewall between **company's net** and **Internet**
  - Firewall between **secret future product group's LAN** and **rest of company's net**
  - Firewall A between **Internet** and **public servers**, firewall B between **servers** and **rest of company's net**
  - Software personal firewall on **desktop machine**
- Filtering policy depends on **which attacks want to defend against**, e.g.,
  - Packet filtering router
  - Application-level gateway (proxy for ftp, HTTP, &c.)
  - Personal firewall disallows Internet Explorer from making outbound SMTP connections

# Background: Internet Services and Port Numbers

- Recall that UDP and TCP protocols identify service by **destination 16-bit port number**
- Well-known services: typically listen on **ports  $\leq 600$** 
  - UNIX: must be root to listen on or send from port  $< 1024$
- Outgoing connections typically use **high source port numbers**
  - App can ask OS to pick unused port number
- See **`/etc/services`** on UNIX host for list of well-known ports

# Non-Secure Services

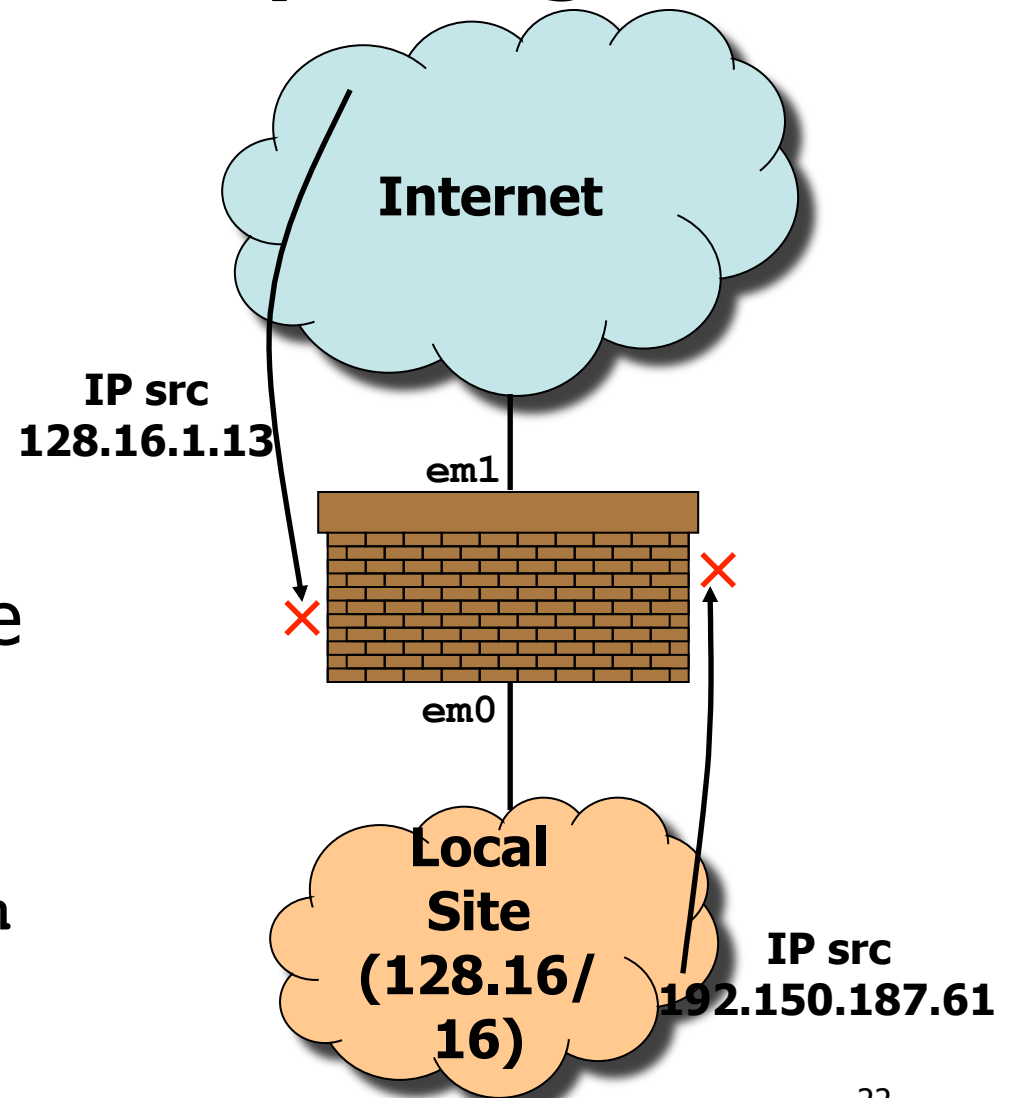
- NFS server (port 2049)
  - Recall: can read/write entire file system given file handle for any directory
  - File handles guessable on many platforms
- Portmap (port 111)
  - Relays RPC requests, so they appear to come from localhost
- FTP (port 21)
  - Client instructs server to connect to self; can instead direct server to connect to 3<sup>rd</sup> party (“bounce” attack)
- Yellow pages/NIS
  - Allows remote retrieval of password database
- Any server with a vulnerability
  - MS SQL (UDP 1434), DNS (53), rlogin (513), lpd (515), ...

# Firewalls: Packet Filtering

- Examine protocol fields of individual packets; filter according to rules
  - IP source, destination addresses
  - IP protocol ID
  - TCP/UDP source, destination ports
  - TCP packet flags (e.g., SYN, FIN, ...)
  - ICMP message type
- Example: to prevent remote lpd exploit, block all inbound TCP packets to destination port 515
  - Remote users shouldn't be printing at your site anyway

# Firewall Example: Blocking Source Spoofing

- Block traffic from outside your site with a source address in your site's address block
- Egress filtering: block traffic from within your site with a source address not in your site's address block
  - e.g., rule:  
`"deny ip not from 128.16/16 recv em0 xmit em1"`



# Firewall Example: Blocking Outbound Mail

- Worms often use infected hosts to send spam or confidential documents
- Defense: authorize only a few servers at site to send outbound mail; filter all outbound mail connections from others
- e.g., rules:

```
allow tcp from 128.16.1.20 not to  
128.16/16 dst-port 25
```

```
deny tcp from 128.16/16 not to  
128.16/16 dst-port 25
```

# **Firewall Example: Block All Inbound Traffic by Default**

- Little control over what software users run on desktops (including servers) at most sites
- May wish to avoid remote exploits of any software run on users' desktops
- Policy:
  - disallow all inbound TCP connections but those to known legitimate servers (e.g., one public web server, one mail server)
  - allow all outbound TCP connections
- Implementation:
  - Stateless way: drop all inbound TCP packets with SYN flag set, but not ACK flag



# Stateful Firewalling

- Stateful way to implement “outbound TCP only”:
  - Firewall stores state for every active TCP connection (src IP, src port, dst IP, dst port)
  - Only forwards “legal” packets for current state
    - e.g., if connection unknown, only allow outbound packets with SYN flag set, but not ACK flag
    - e.g., if connection known, only allow inbound packets with data after SYN/ACK seen
  - Time out connection state for long-idle connections
- Also used to block inbound UDP only
  - No standard SYN, ACK fields in UDP to support stateless filtering
- Risk: state memory exhaustion on firewall

# Firewalling Complex Protocols

- Consider FTP
- Client connects to server, instructs server to open TCP connection back to client on specified client-side port
- Client's firewall won't allow inbound connection!
- One solution: application-level proxy
  - Client's firewall starts FTP application-level proxy upon detecting FTP session
  - Proxy on firewall acts as client for TCP connections with remote server, server for TCP connections with local client
  - Can enforce policy for many protocols (SMTP, HTTP, &c.)
  - But not used for encrypted protocols (SSL, SSH, &c.)